# TOWARDS AFFORDABLE EXTERNALLY CONSISTENT GUARANTEES FOR GEOREPLICATED SYSTEMS

Manuel Bravo Luís Rodrigues





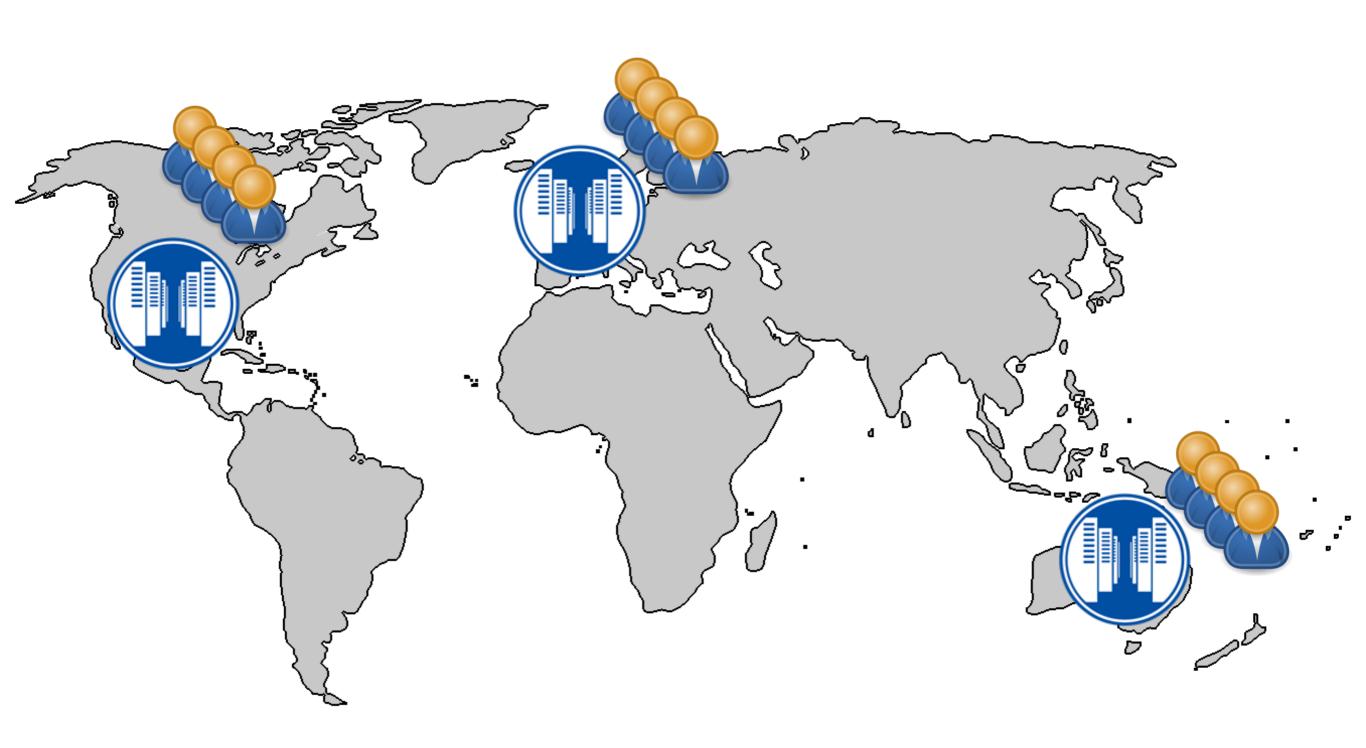




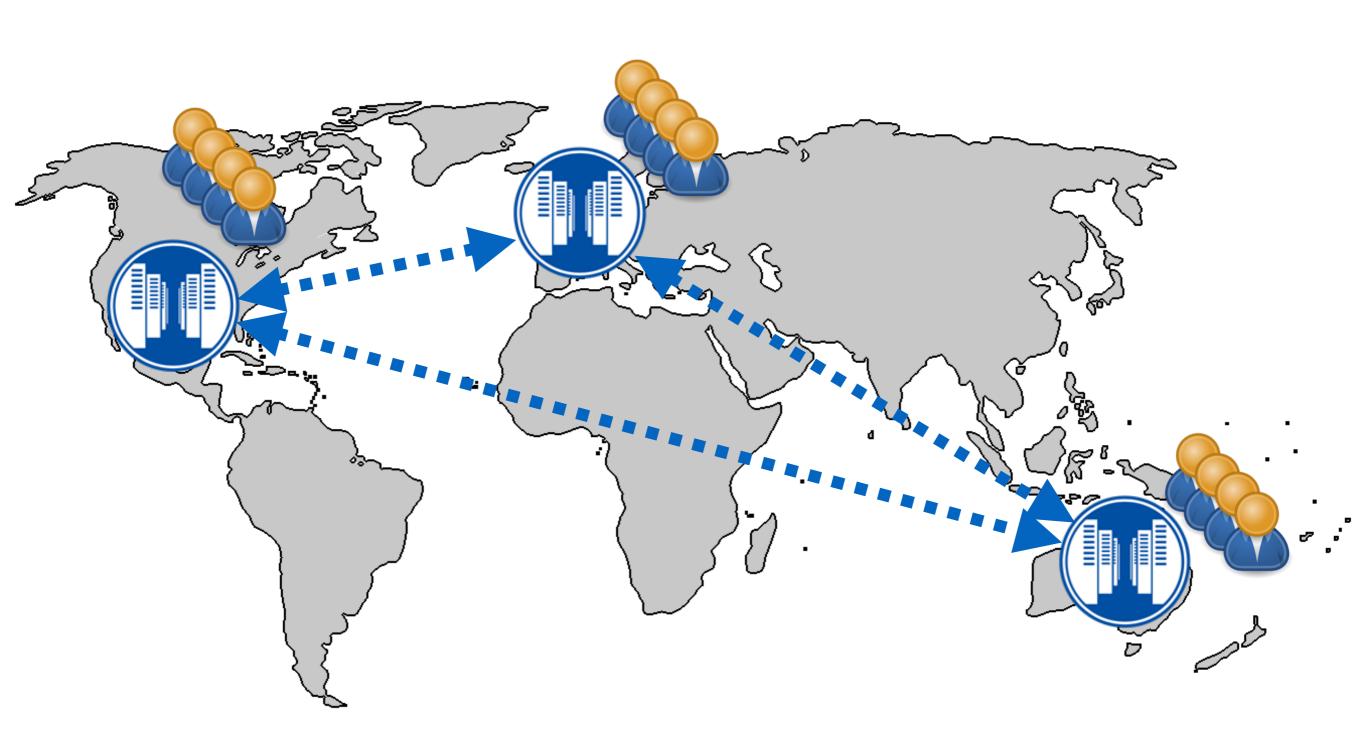
#### Geo-replication



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strong eventual consistency

strong consistency eventual consistency



#### natural semantics

#### high latency

e.g., Google Spanner [OSDI'12]

strong consistency eventual consistency



#### natural semantics

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e.g., Google Spanner [OSDI'12]

#### **low latency**

#### error-prone

e.g., Dynamo [SOSP'07]

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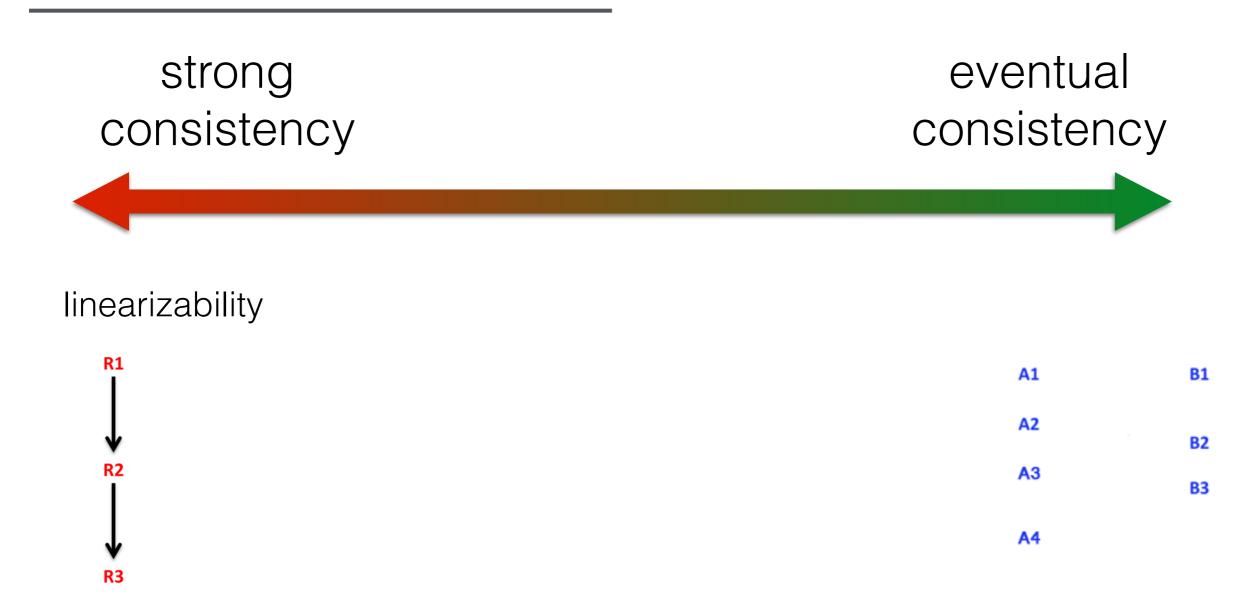
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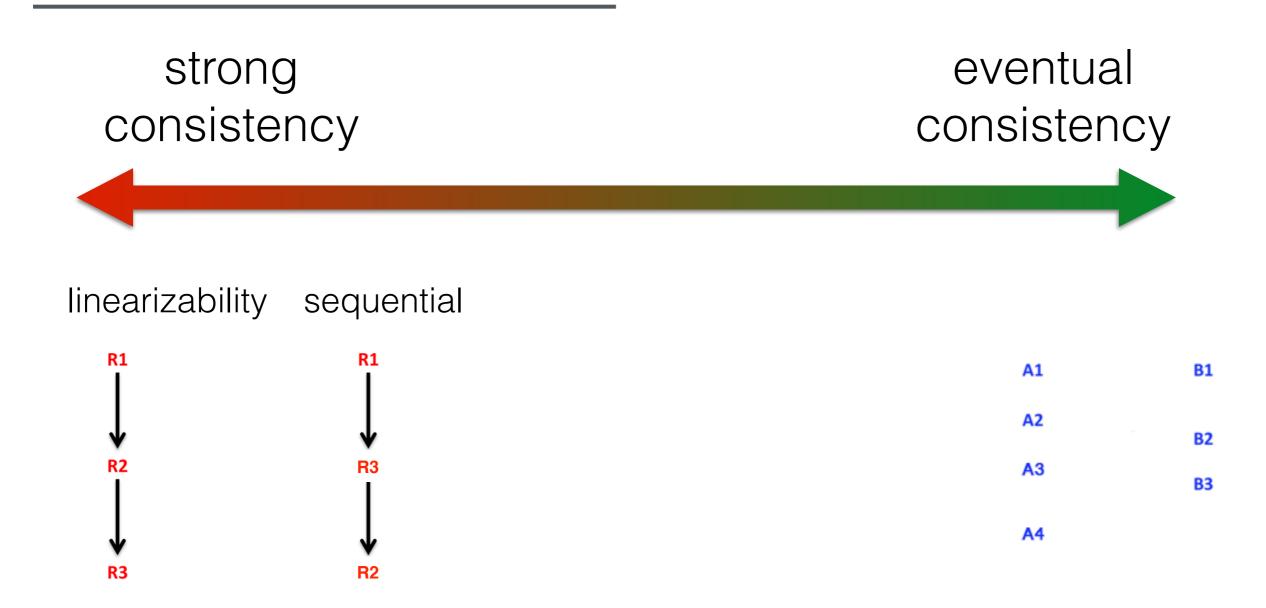
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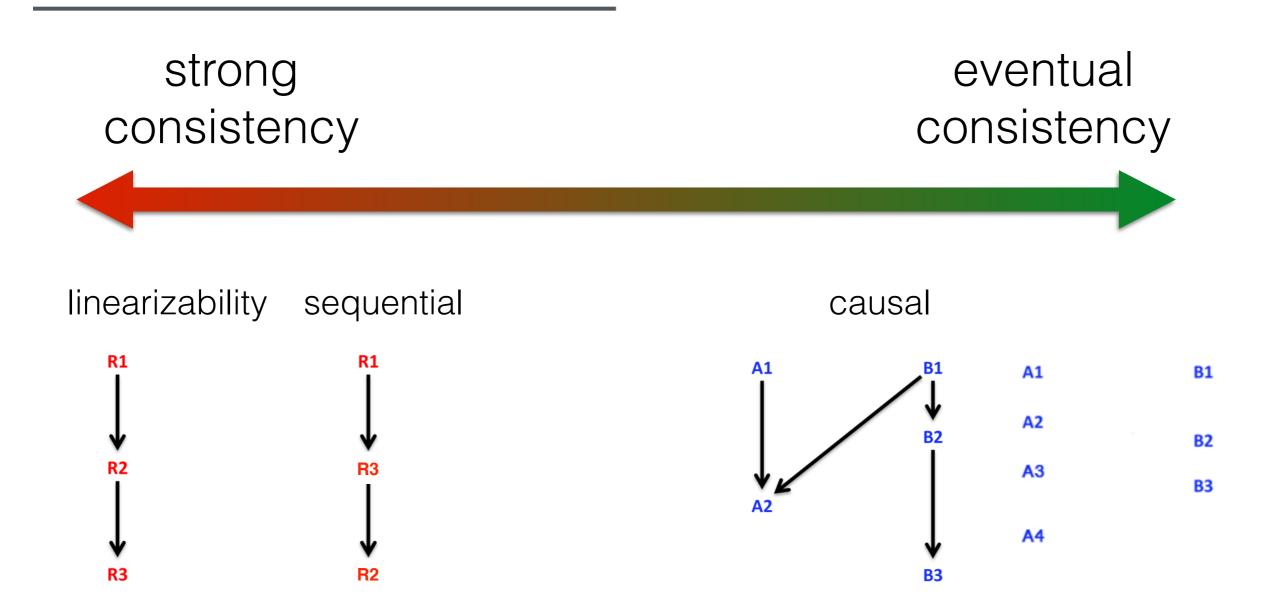


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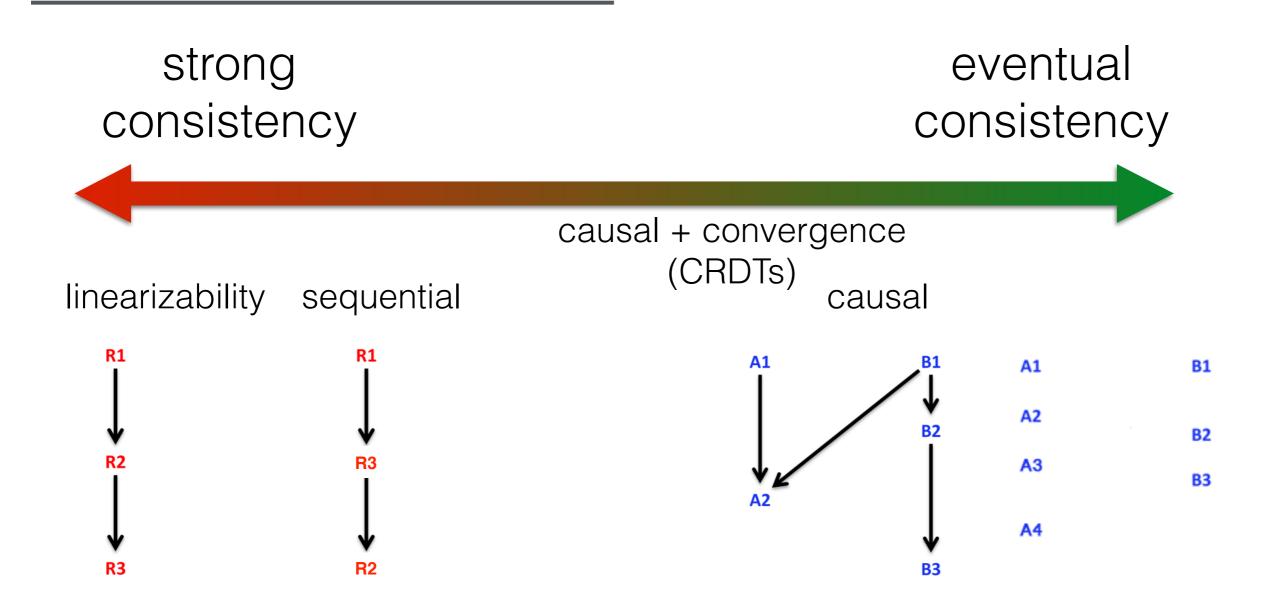


natural semantics
high latency



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#### Alternative

### mixing semantics, consistent when necessary

**Parallel Snapshot Isolation** 

[SOSP'11]

**RedBlue Consistency** 

[OSDI'12]

**Explicit Consistency** 

[EuroSys'15]

**Pileus** 

[SOSP'13]

**Session guarantees** 

[SOSP'97]

#### Alternative

mixing semantics, consistent when
Interestingly, most
previous solutions do not solution

RedBlue Coconsider external—or

Blobal guarantees

Pileus
[SOSP'13]

Session guarantees [SOSP'97]

#### External consistency

## clients are served with a view of the system consistent to how an external observer would witness the succession of events

#### **CHAPTER 3: TRANSACTIONAL STORAGE**

consider read and write actions, and we will assume that a transaction performs at most one read and one write action on a data element.

The question of serial consistency arises when a transactional storage system concurrently executes actions drawn from several transactions. Imagine two simple transactions:

T1 T2 (a11) Read Y (a21) Read X (a12) Write X (a22) Read Y

(a13) Write Y

The order in which the actions of T1 and T2 are processed is called a *schedule*. A schedule is an arbitrary interleaving of the actions of a set of transactions into a single sequence. A *serial schedule* results when transactions are executed one at a time to completion. Thus, there are two possible serial schedules for our example:

(a23) Write X

Sa: {all, al2, al3, a21, a22, a23} Sb: {a21, a22, a23, al1, al2, al3}.

A schedule generates a *dependency* relation. A dependency relation describes transactions that depend on one another. If S is a schedule, then  $\langle T_n \rangle$ , e.  $T_n \rangle$  is a member

#### very powerful semantics

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an operation observes all operations that were completed as when the operation began

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can cope with back-channeling!

#### very powerful semantics

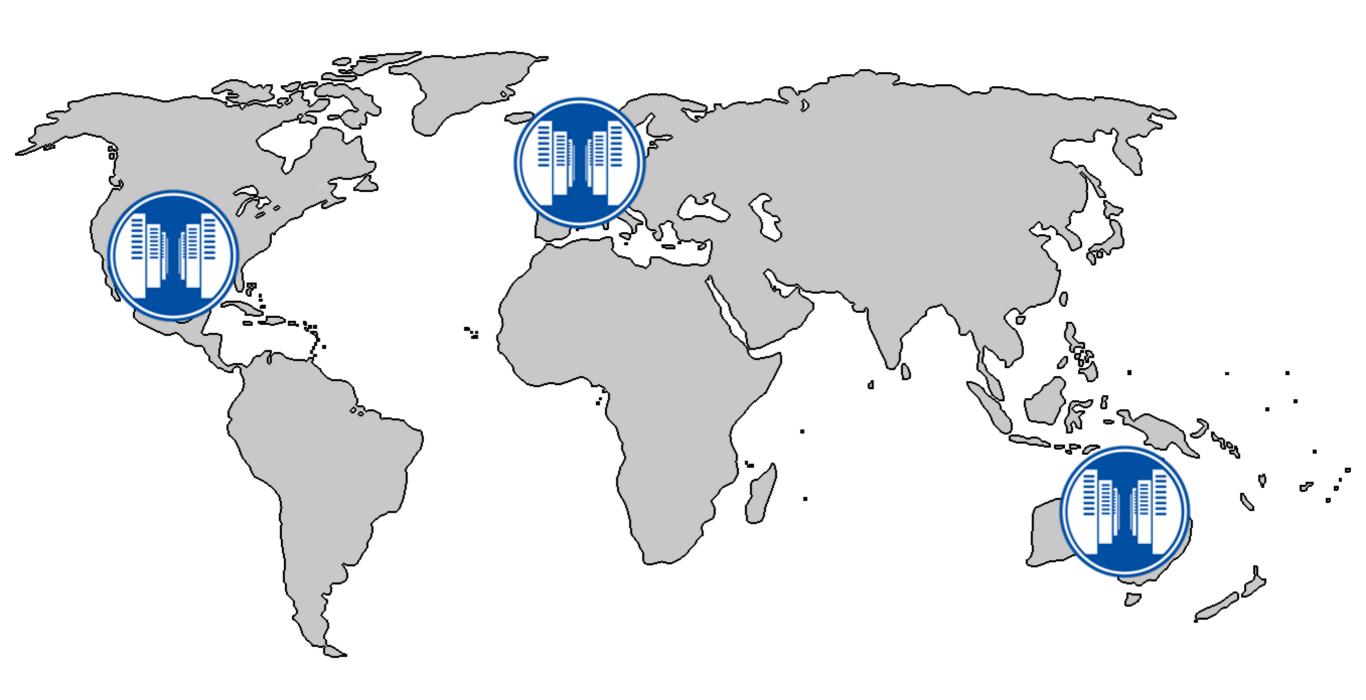
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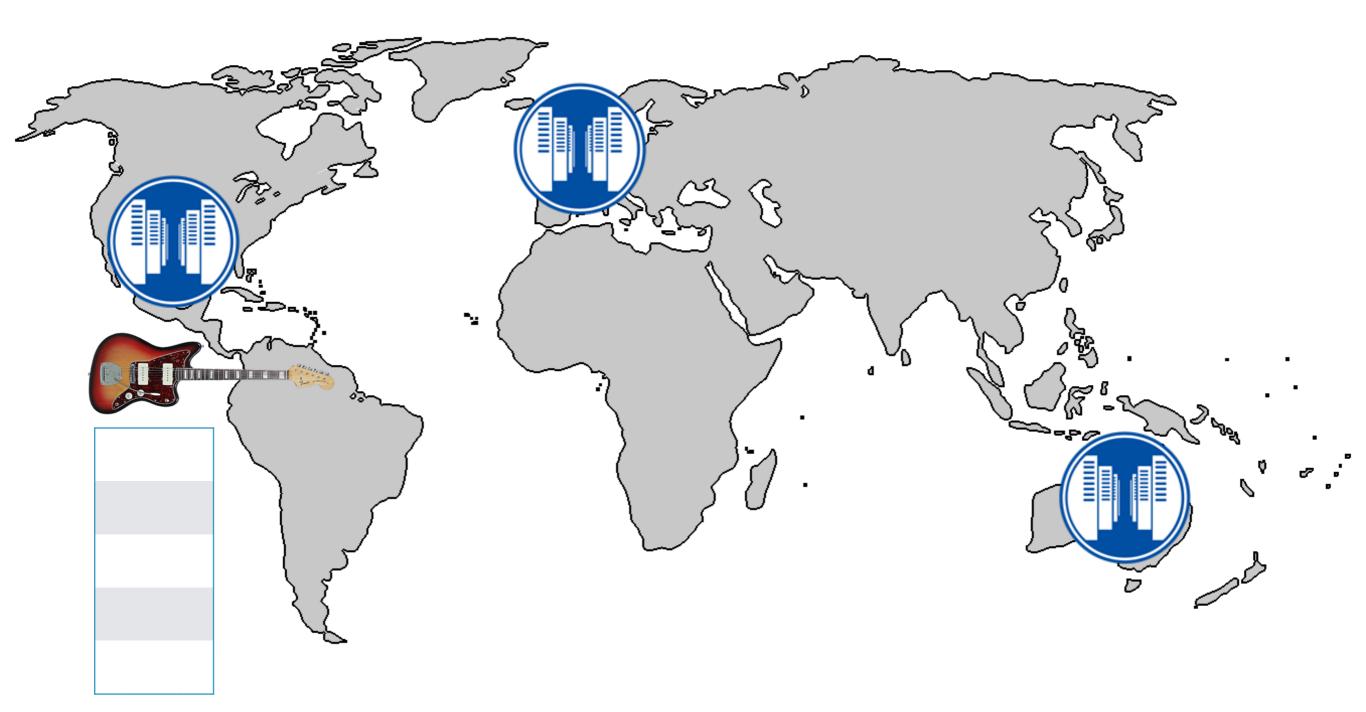
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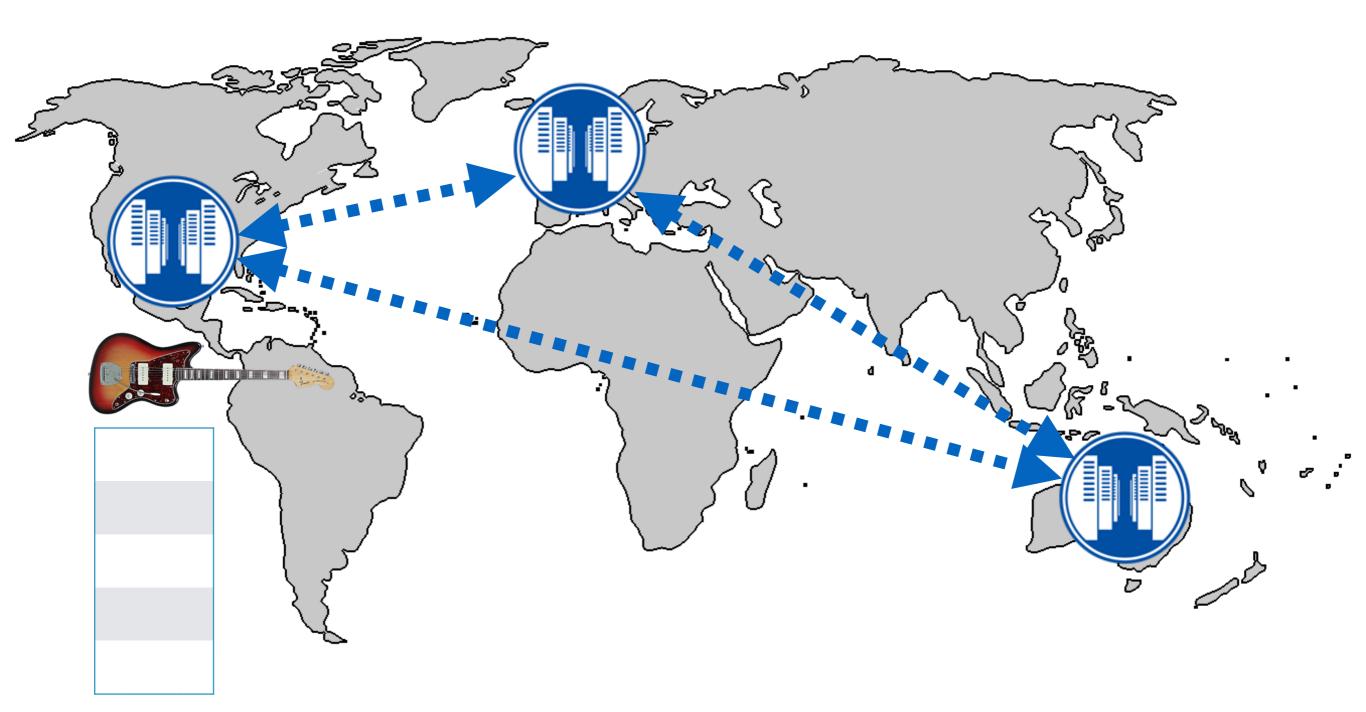
but, requires a lot of coordination then, **very expensive** 

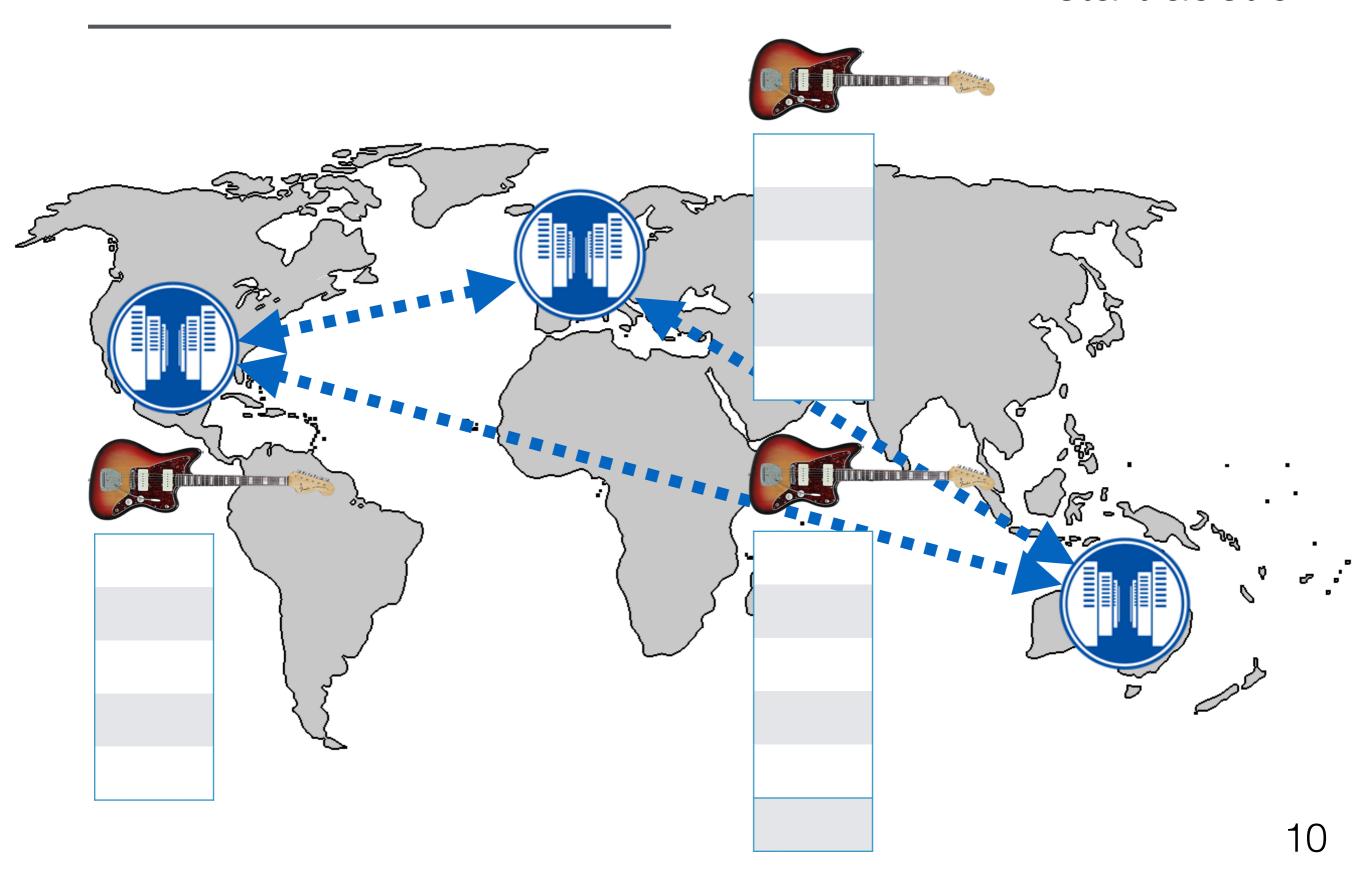
two roles: auctioneers and buyers

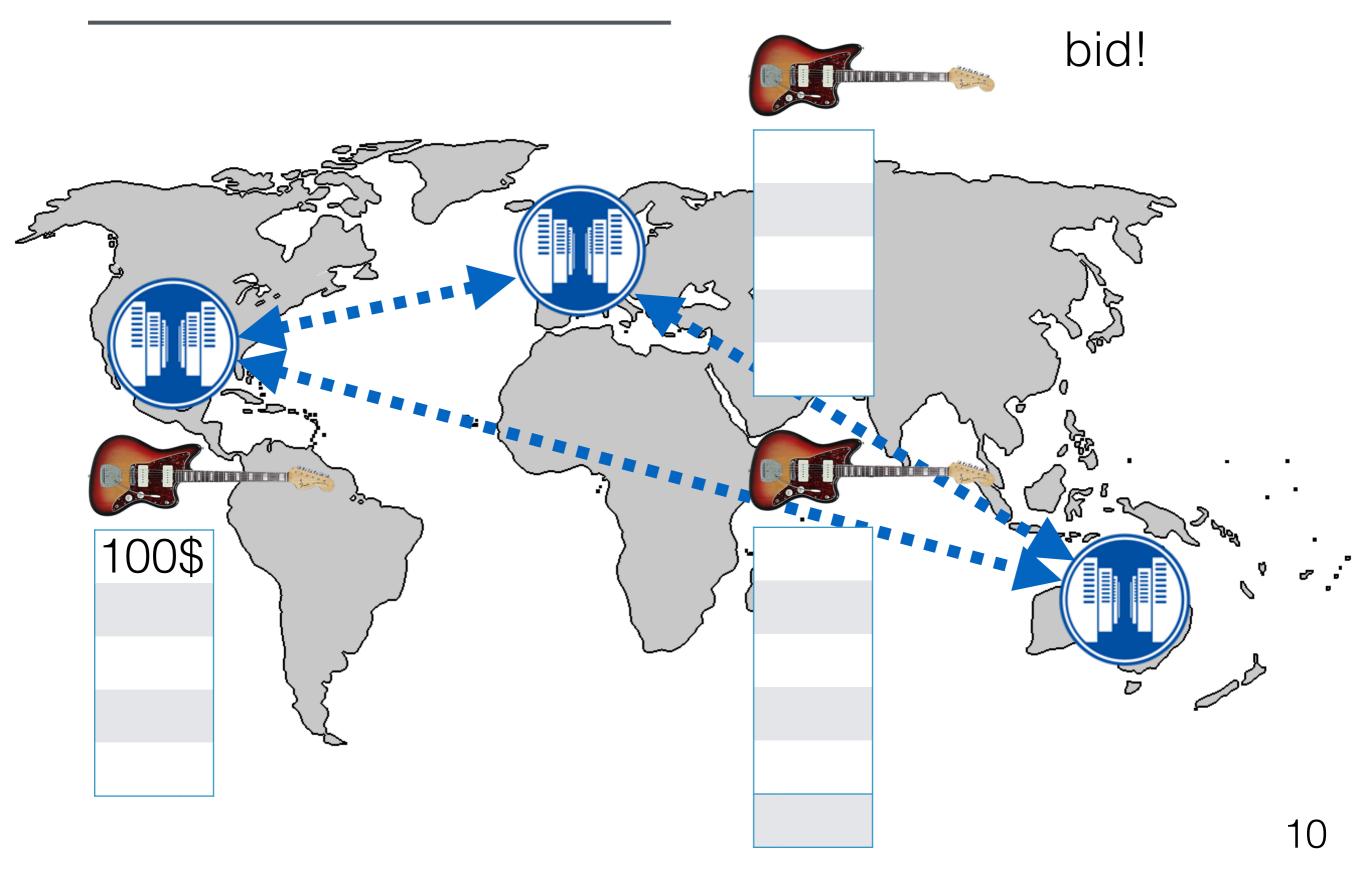
operations (among many): start an auction, place a bid, close an auction, report the winner

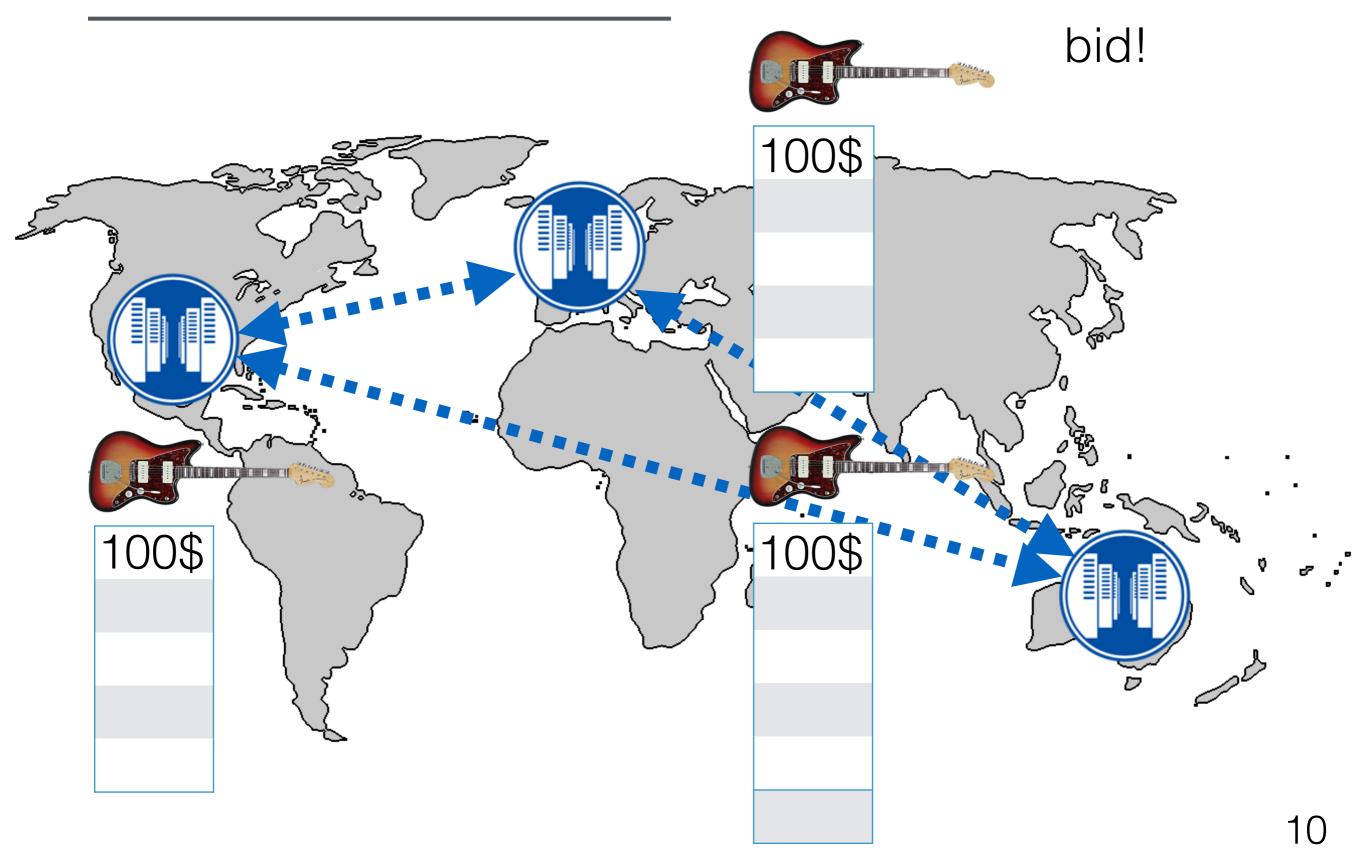




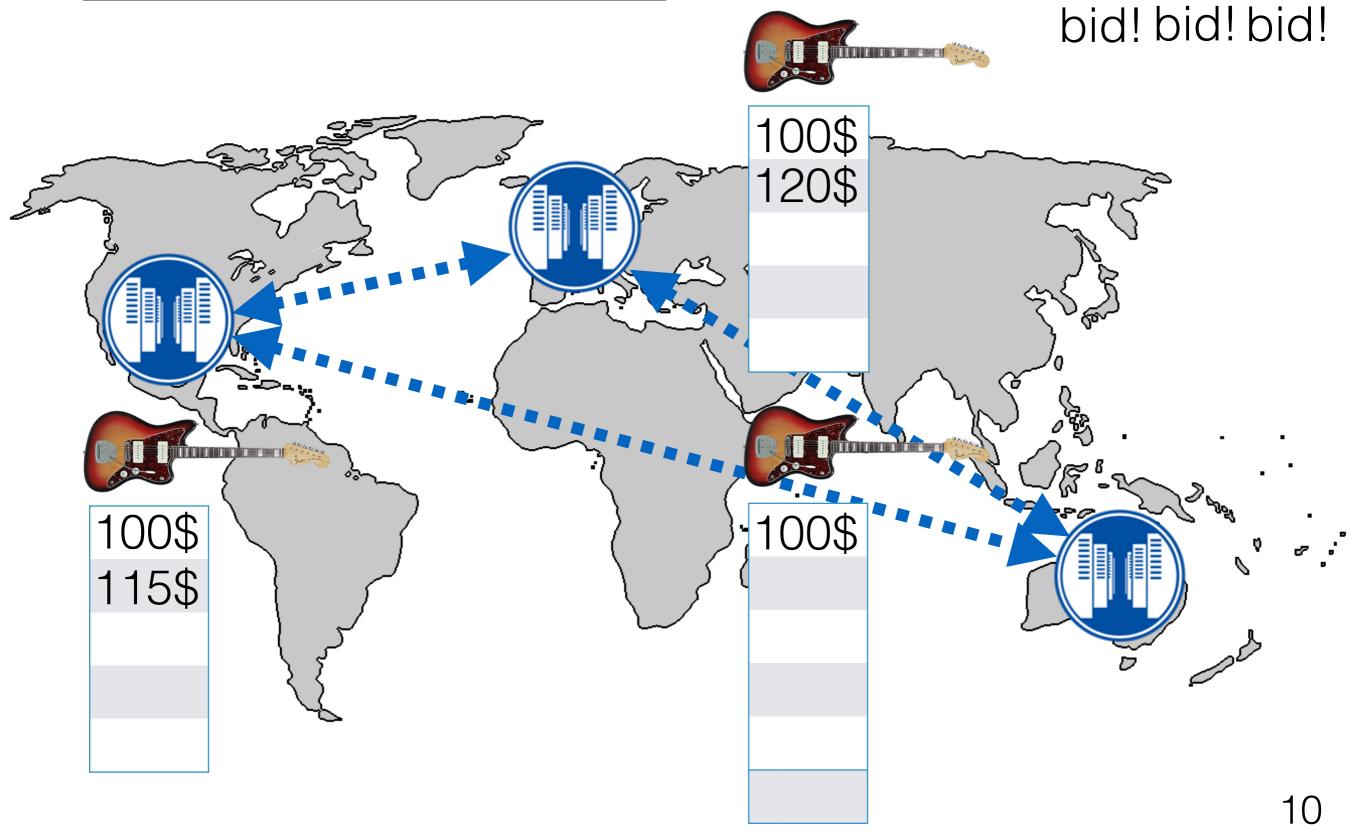


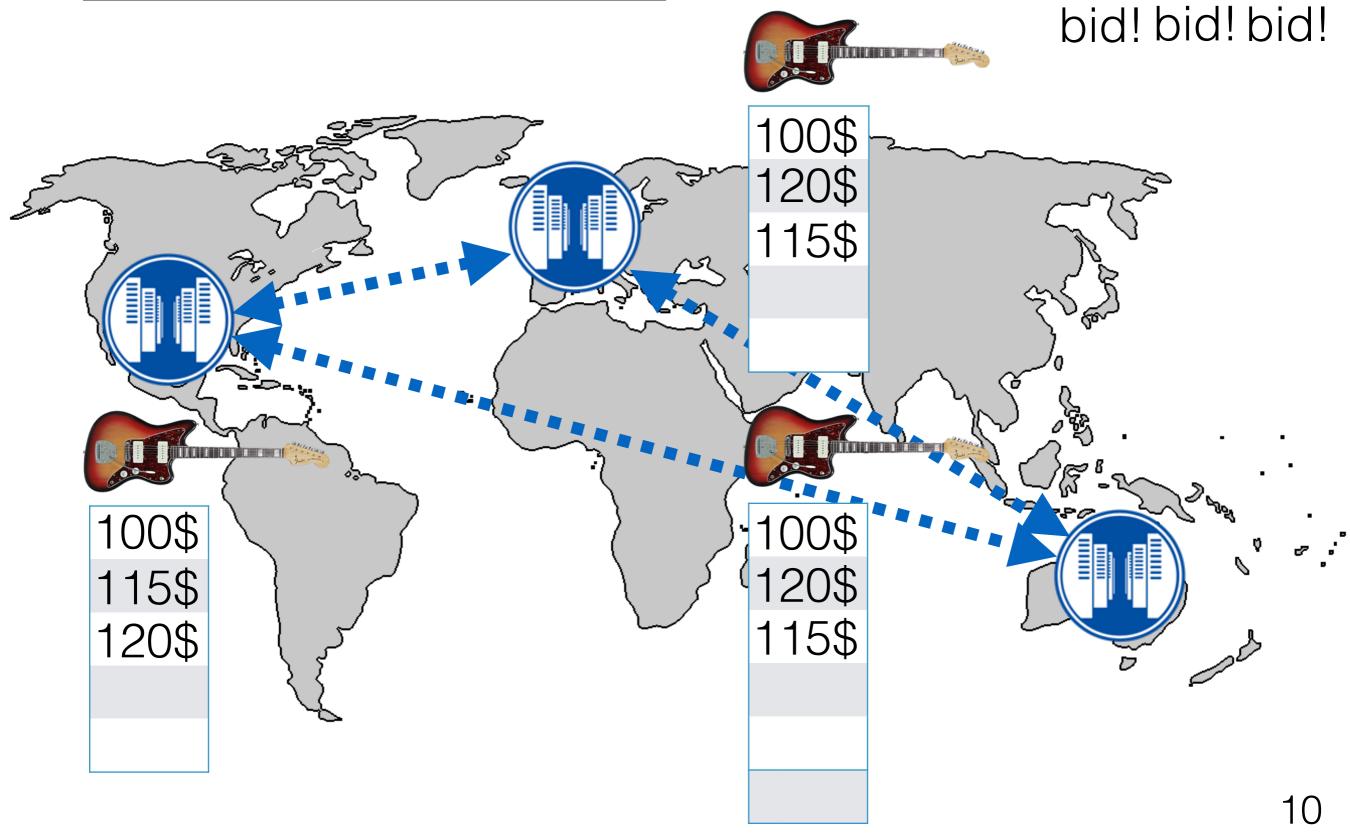






start auction! bid! bid!





#### Use case: an auction service start auction! bid! bid! bid! bid! 100\$ 120\$ 115\$ 100\$ 115\$ 120\$ 115\$ 130\$

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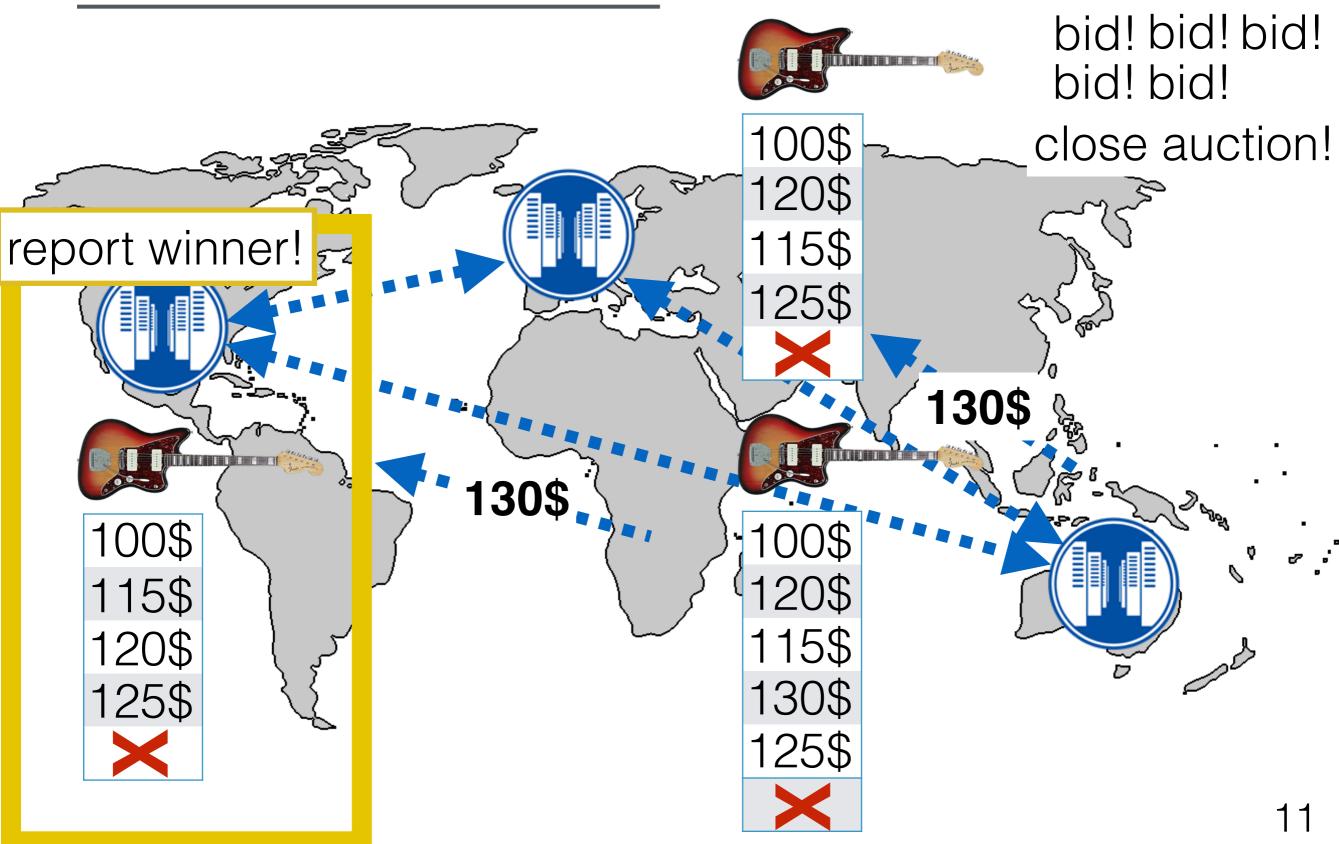
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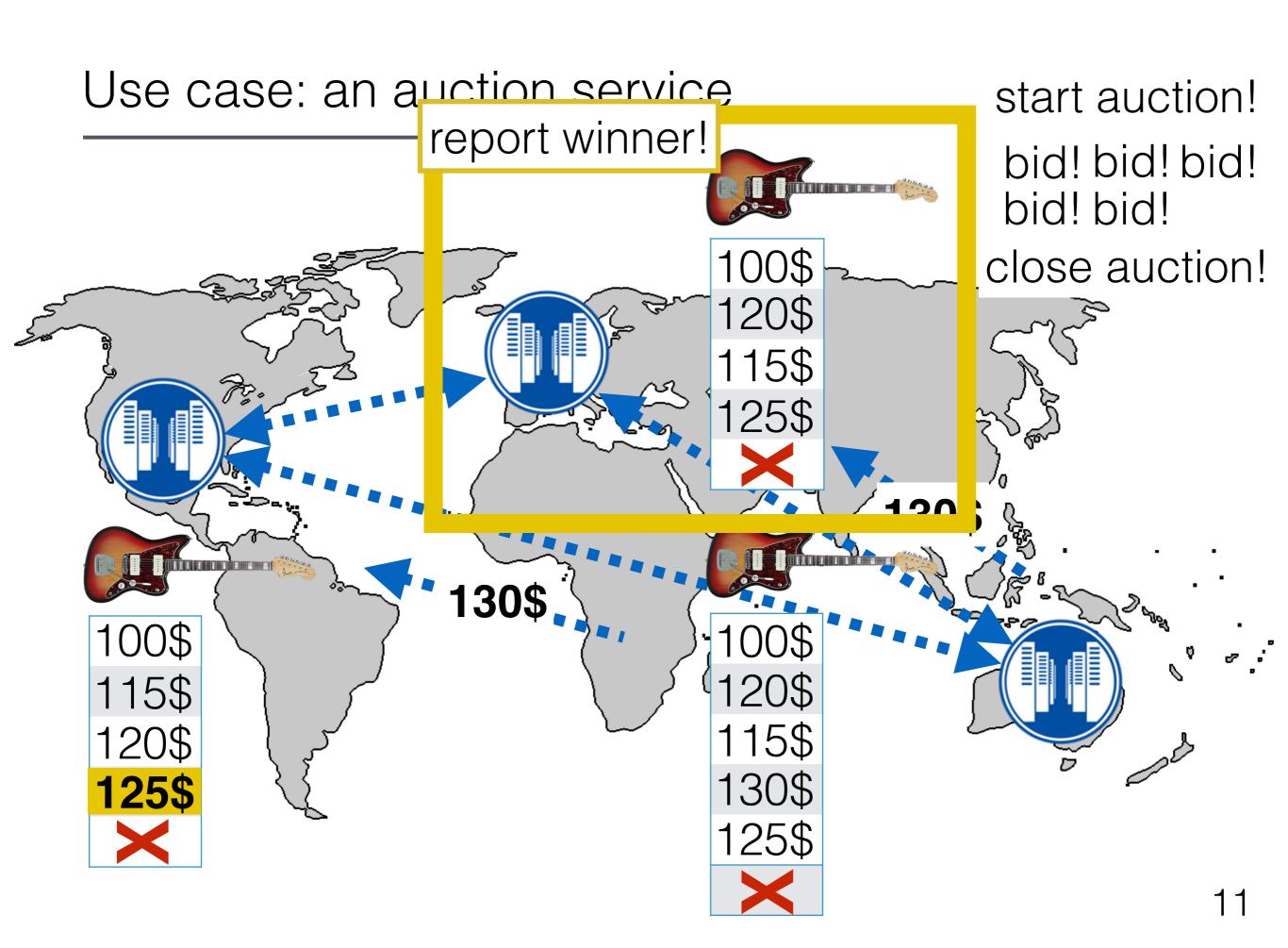
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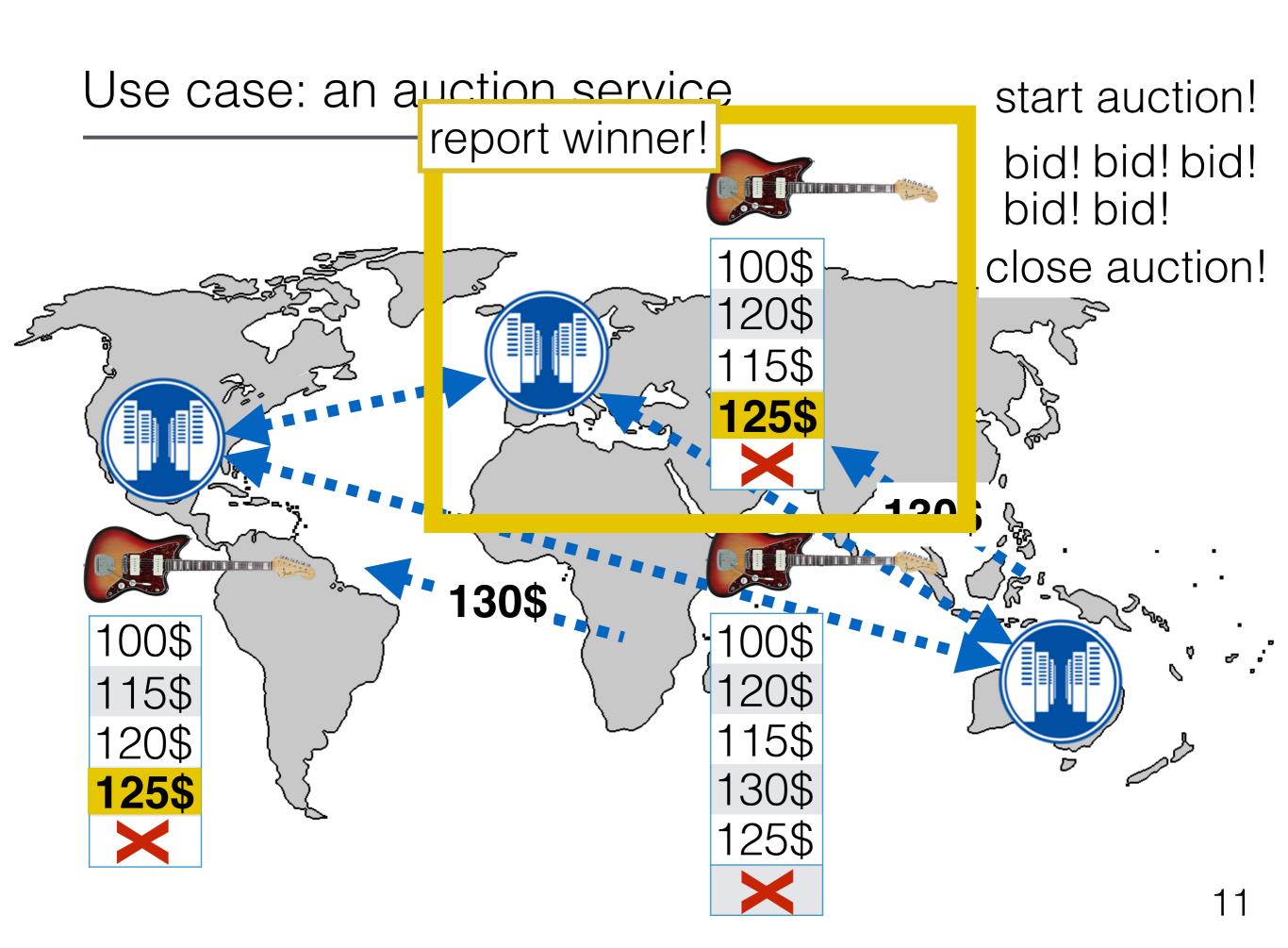
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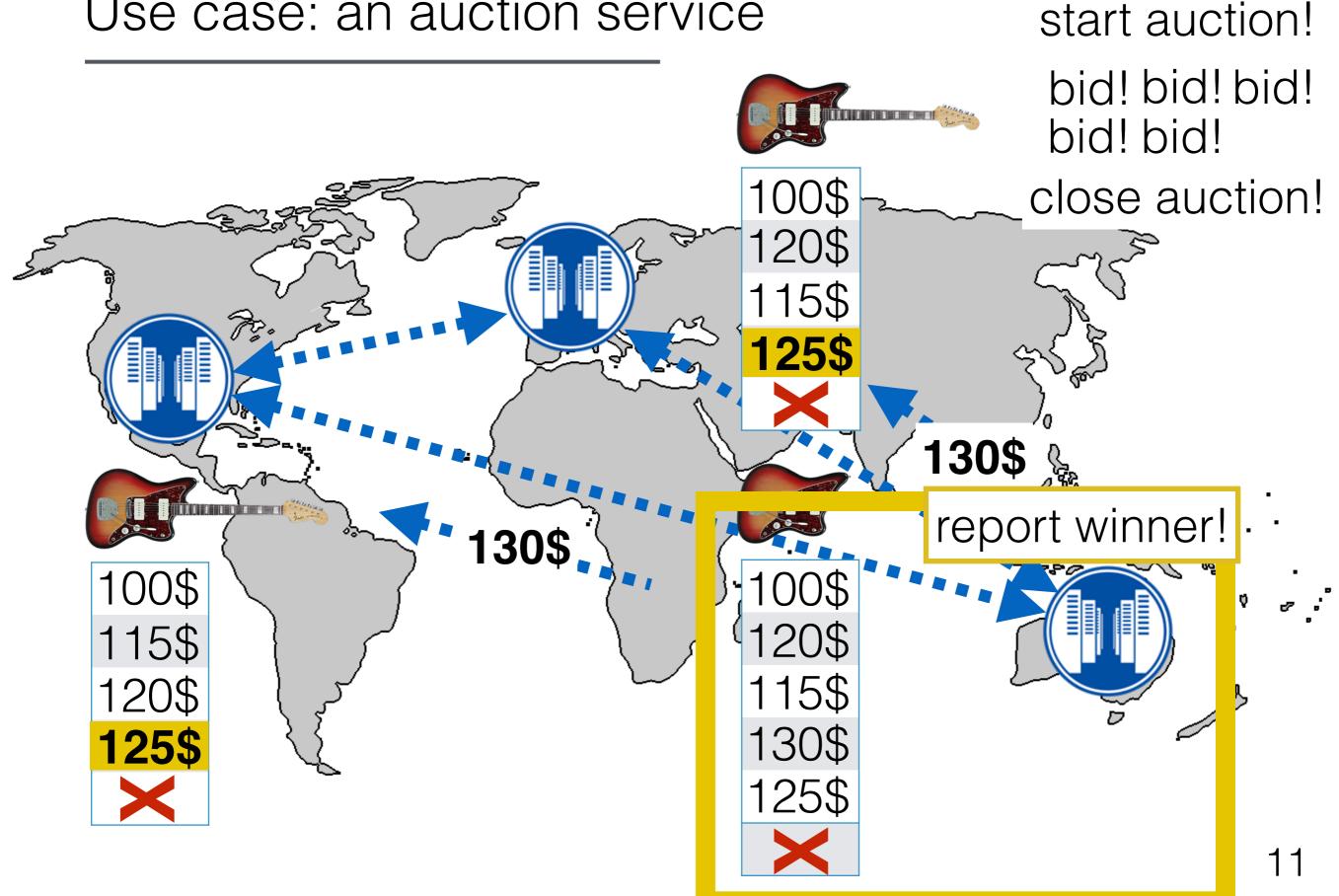
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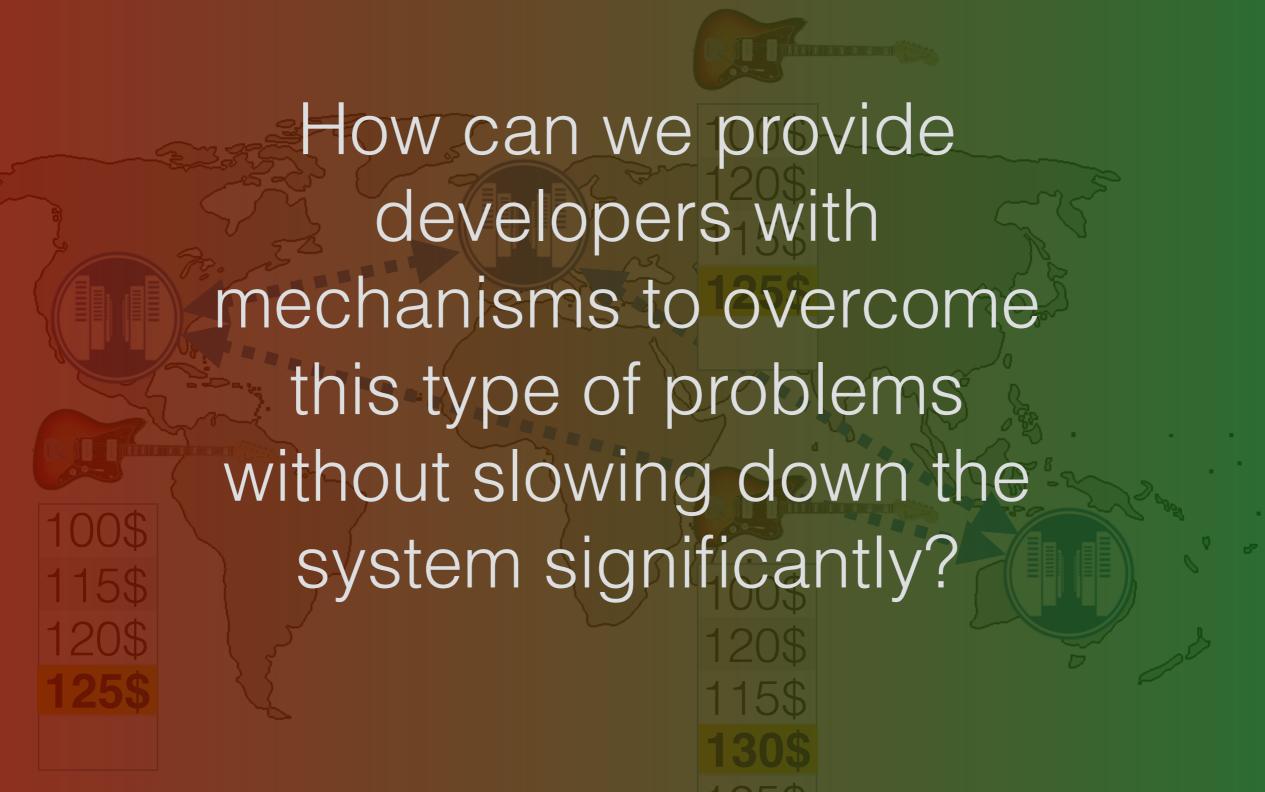
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#### Our work

addresses this inherent tension between performance and meaningful semantics

we present a new consistency model: external causality

takes causal consistency (strongest available) and spice it up with external consistency guarantees

### External causality

internal operations: read from a causally consistent snapshot

external operations: read from a casual snapshot that includes latest updates as of the time when the operation began

### Hypothesis

internal operations are **highly predominant** and **cheap** to implement (and we know how to)

external operations—although expensive—are rarely required

## What about concurrency?

external causality **allows concurrency**, only guaranteeing that operations (both internal an externals) are executed in causal order at each site

this is not ideal from the semantics point of view, but it has significant benefits in practice



internal external causality

Site **a** Site **b** 

Real Time



internal external causality





Real Time



Site **b** 

**b.1** 

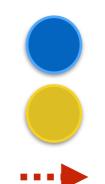
internal external causality

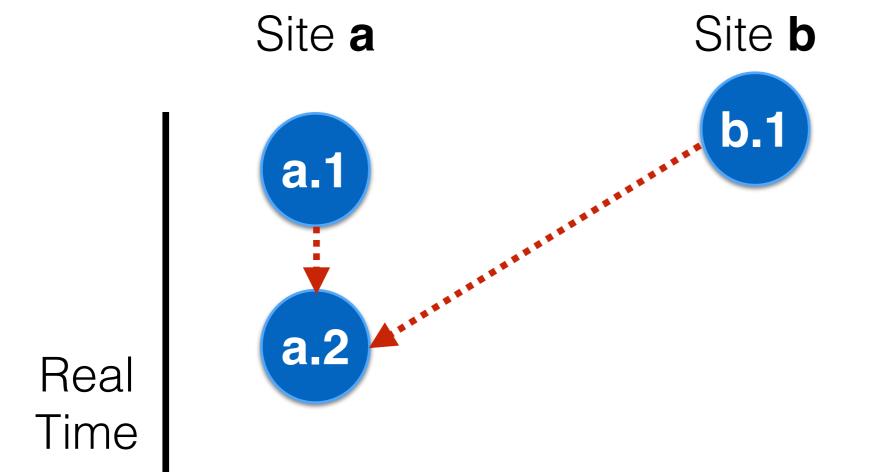
Site a

a.1

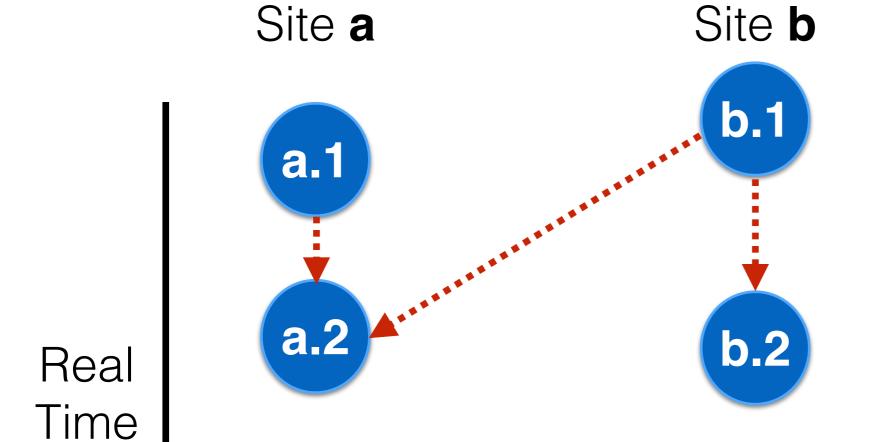
1

Real Time

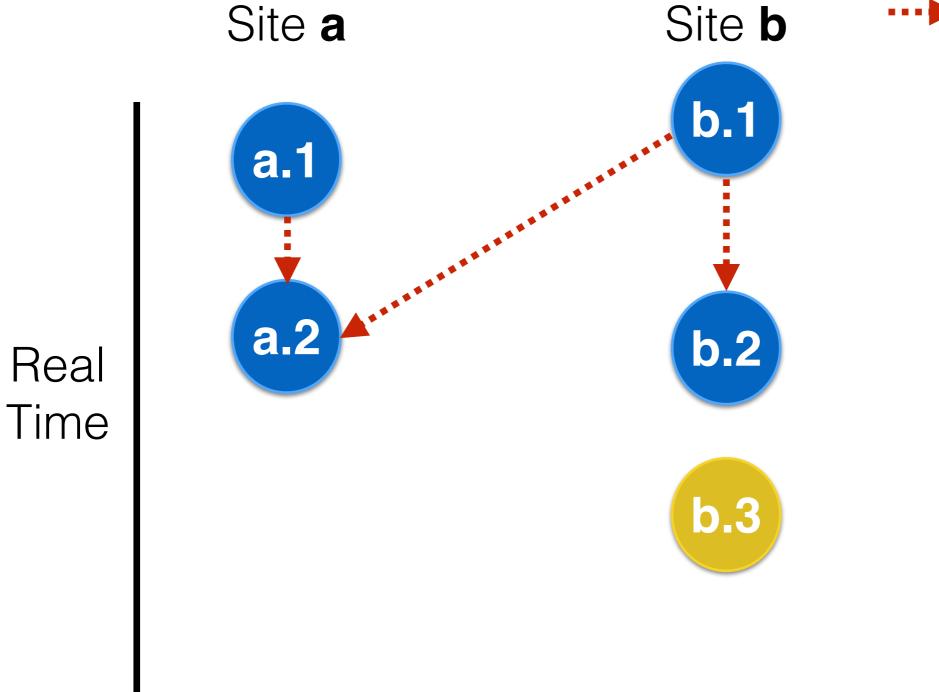






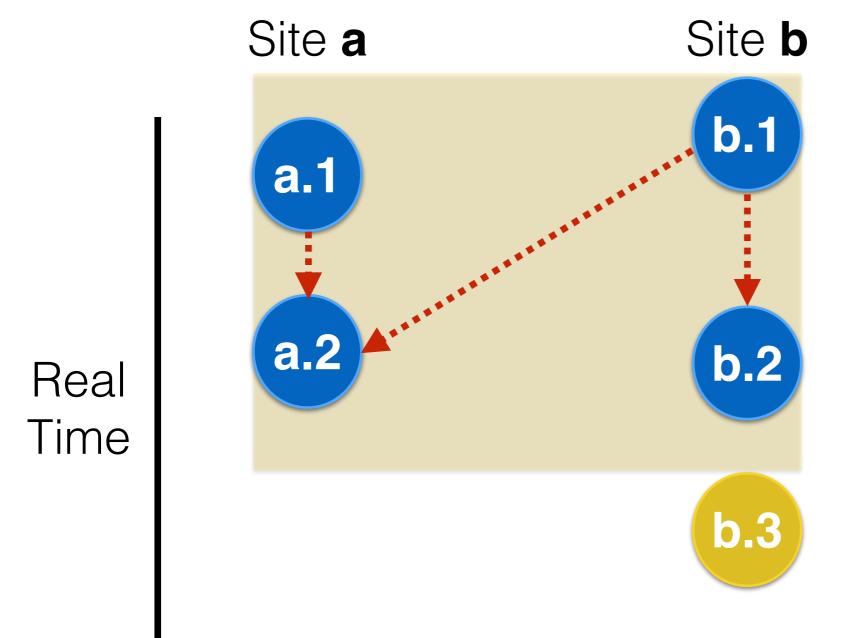








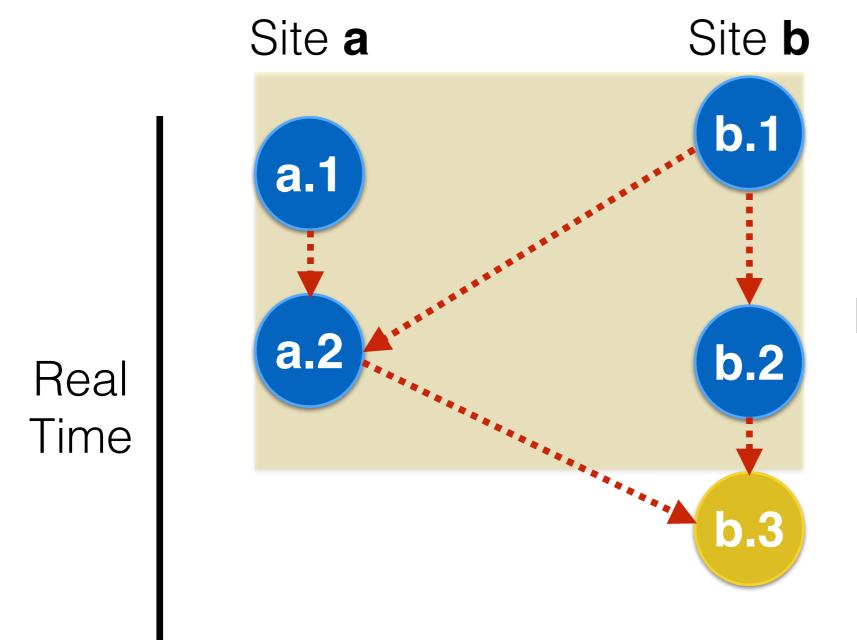
internal external causality



Anything committed locally and elsewhere

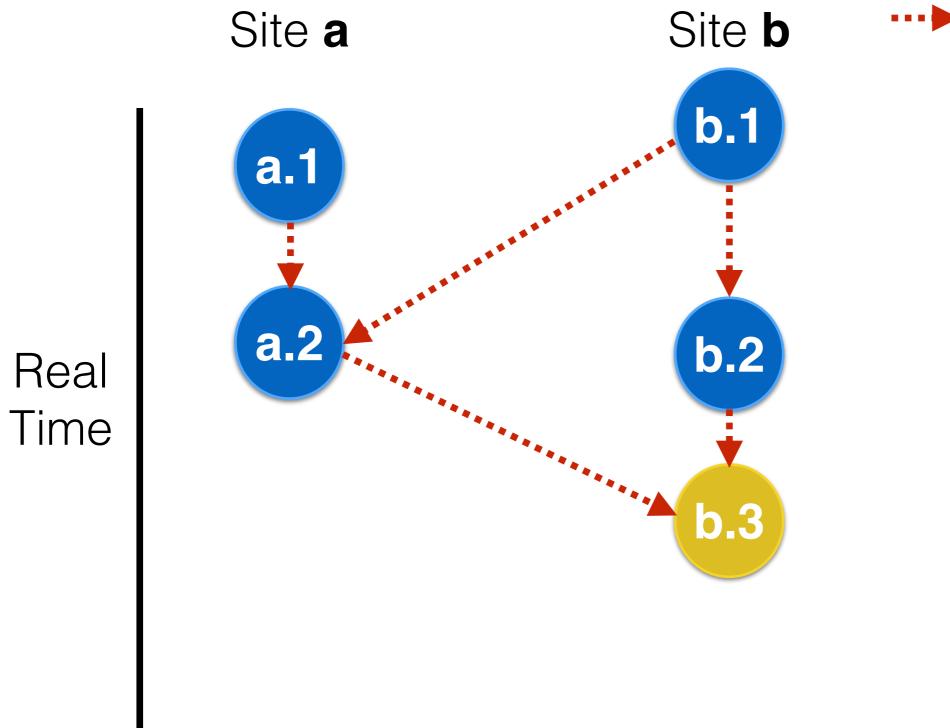


internal external causality

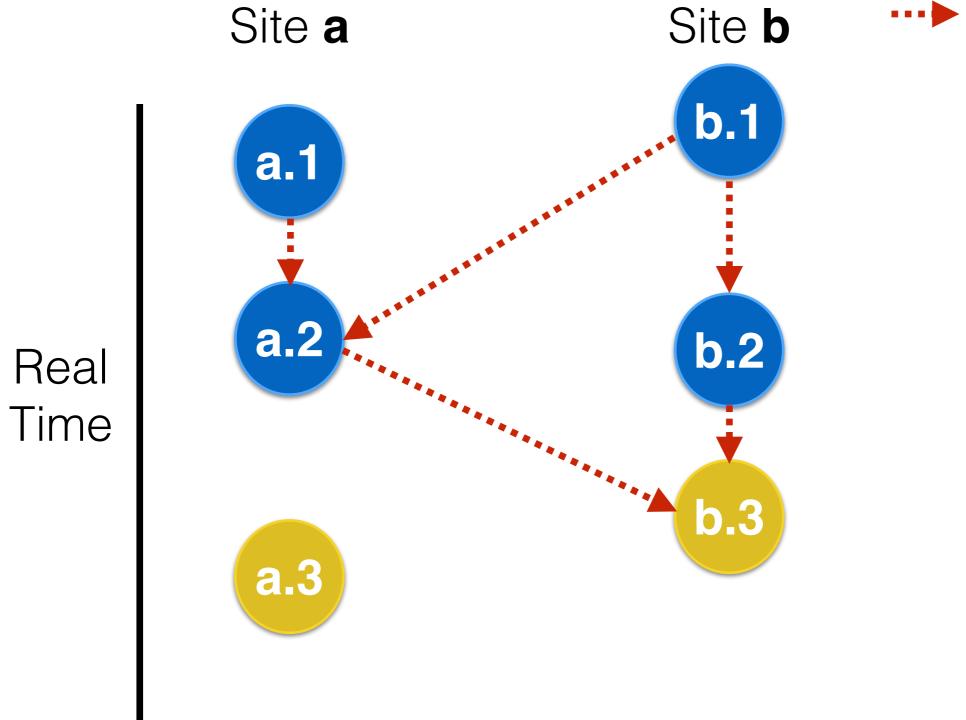


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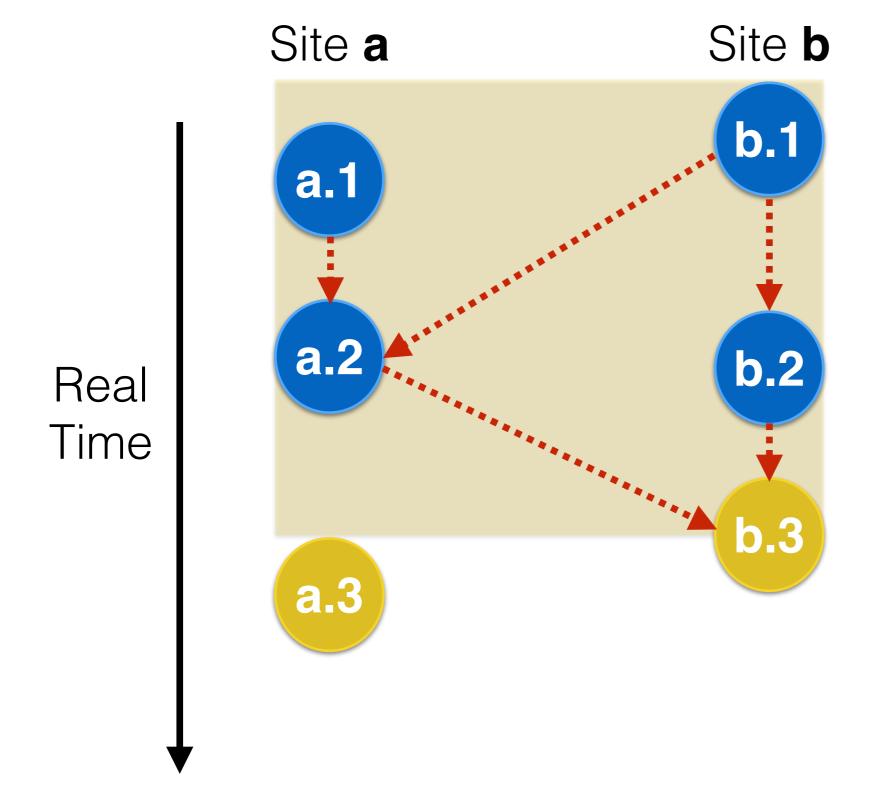




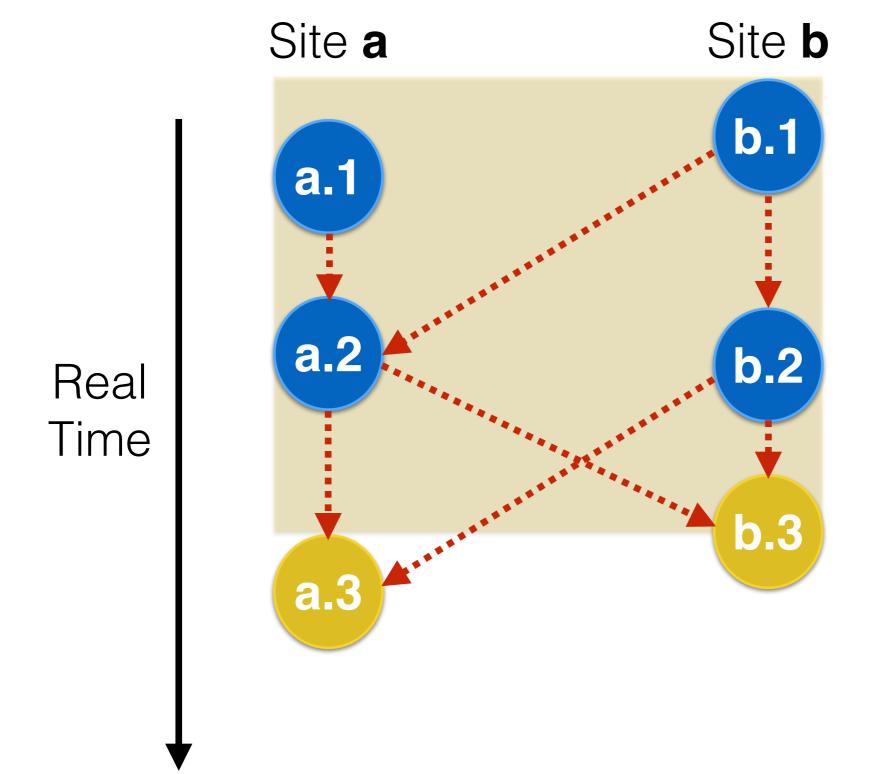




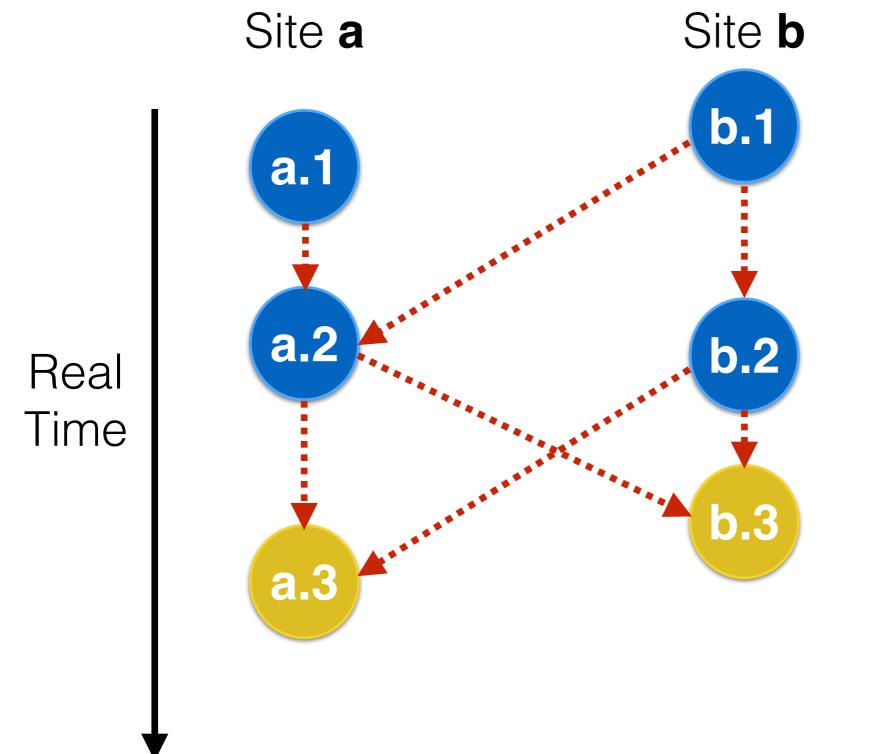




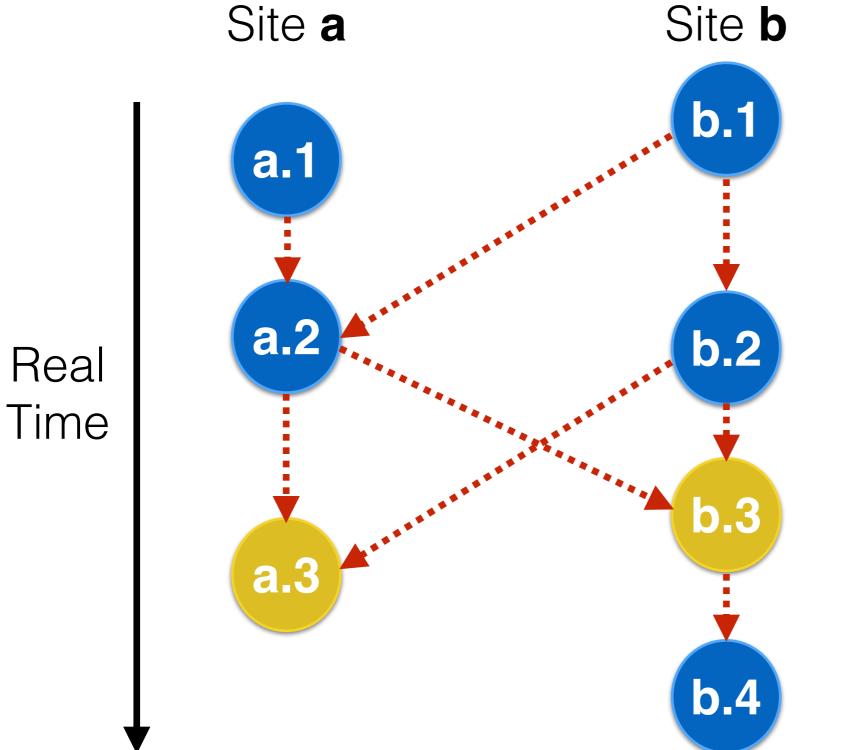




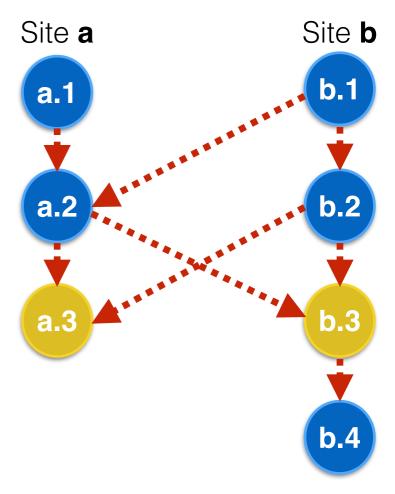




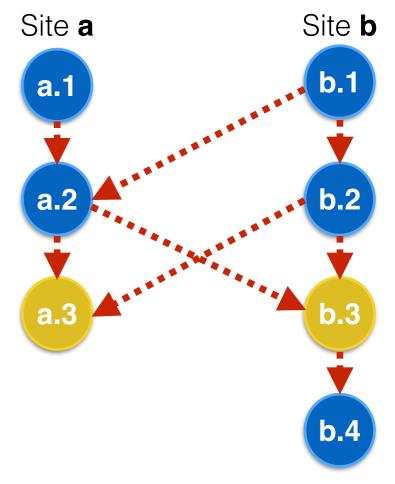




## External causality: causal serialisations of

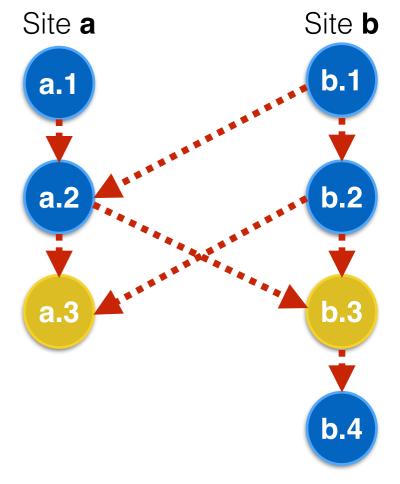


## External causality: causal serialisations of



Site a 
$$a.1 \rightarrow b.1 \rightarrow a.2 \rightarrow b.2 \rightarrow a.3 \rightarrow b.3 \rightarrow b.4$$

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$$a.1 \rightarrow b.1 \rightarrow a.2 \rightarrow b.2 \rightarrow a.3 \rightarrow b.3 \rightarrow b.4$$

Site b 
$$b.1 \rightarrow b.2 \rightarrow a.1 \rightarrow a.2 \rightarrow b.3 \rightarrow b.4 \rightarrow a.3$$

#### Practical issues

internal operations can read from the local site, without prior communication

external operations require prior communication prepare phase—to compute the external snapshot from which they are gonna read from

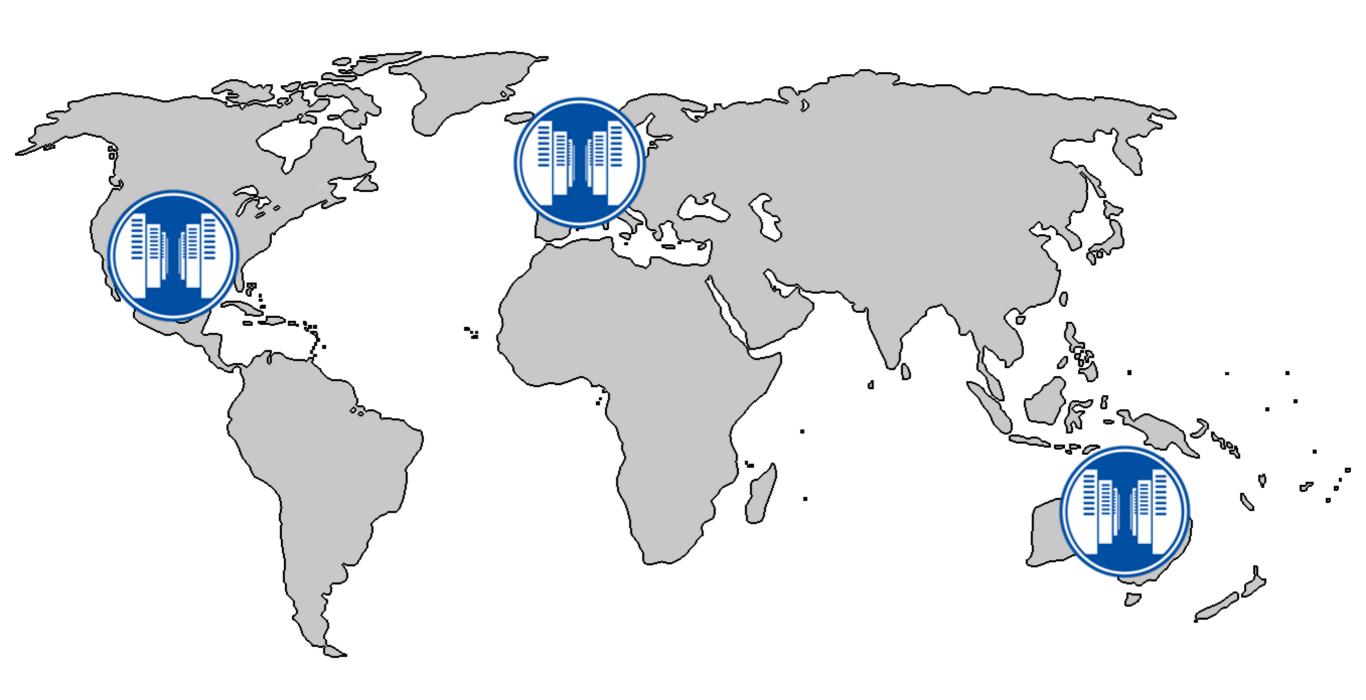
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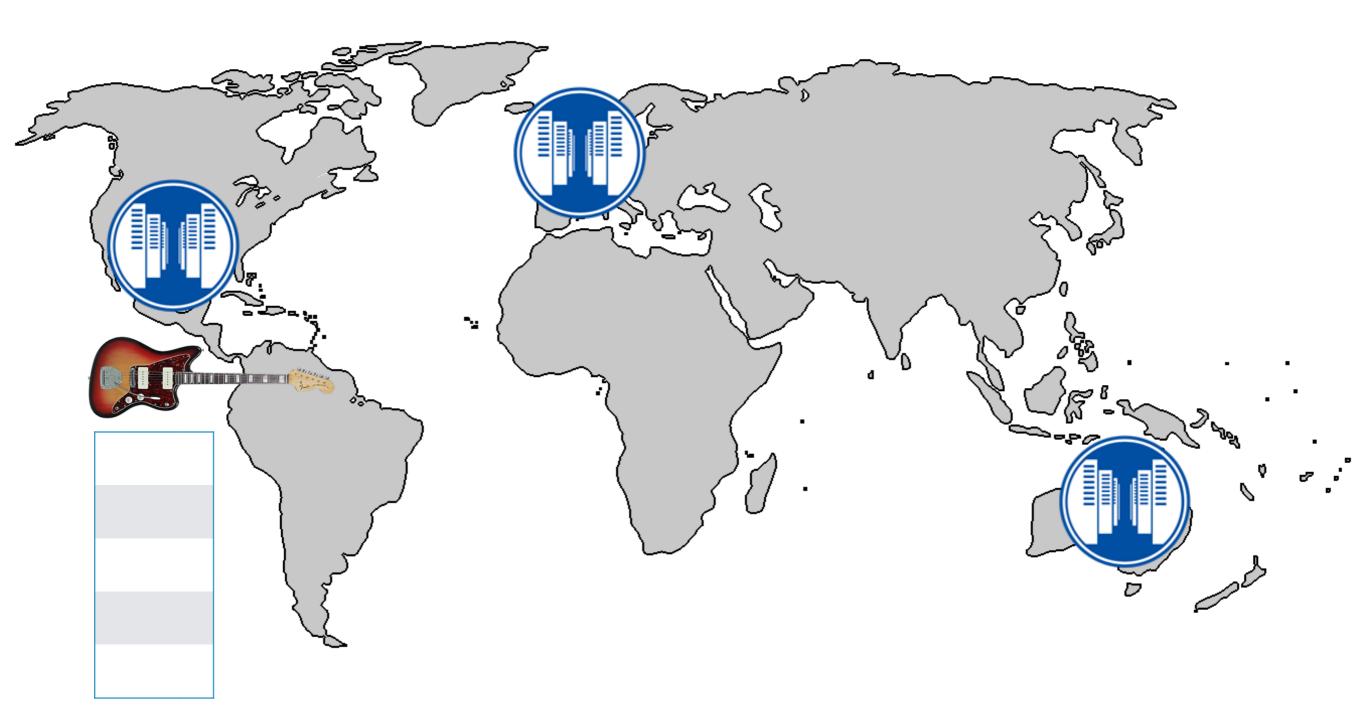
two roles: auctioneers and buyers

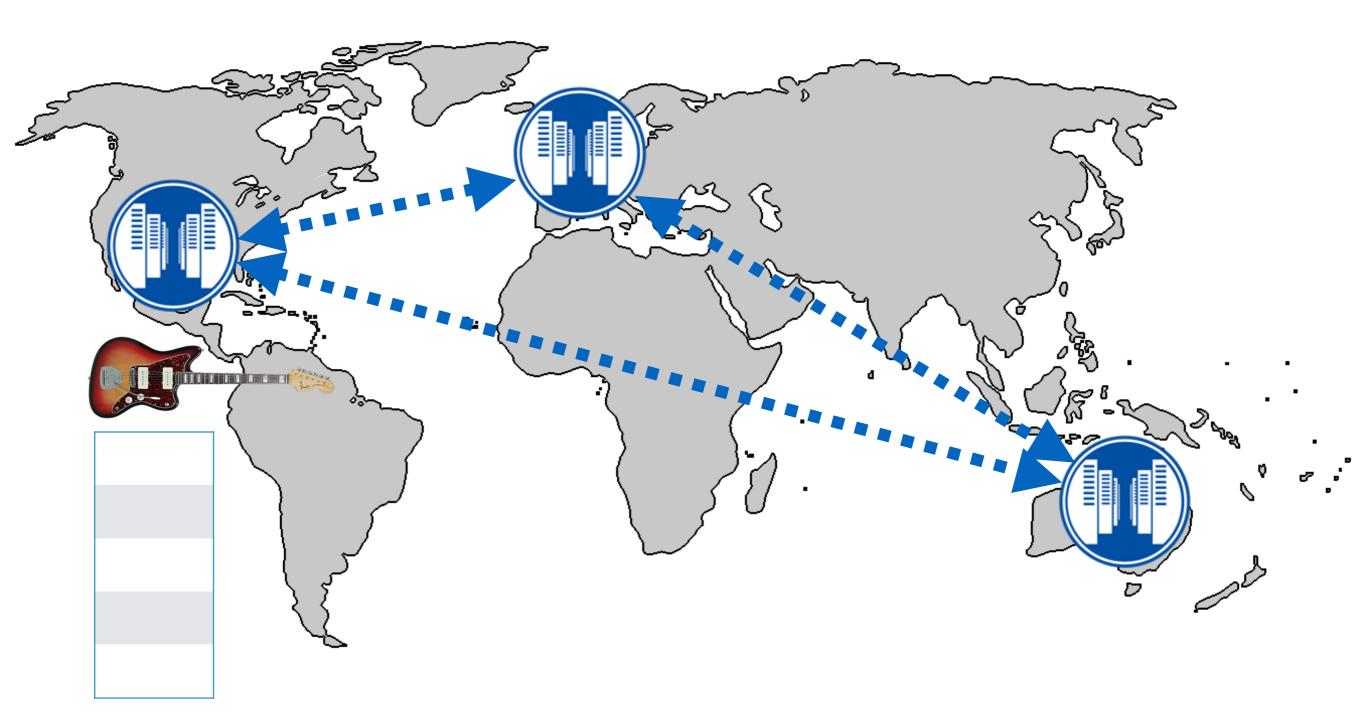
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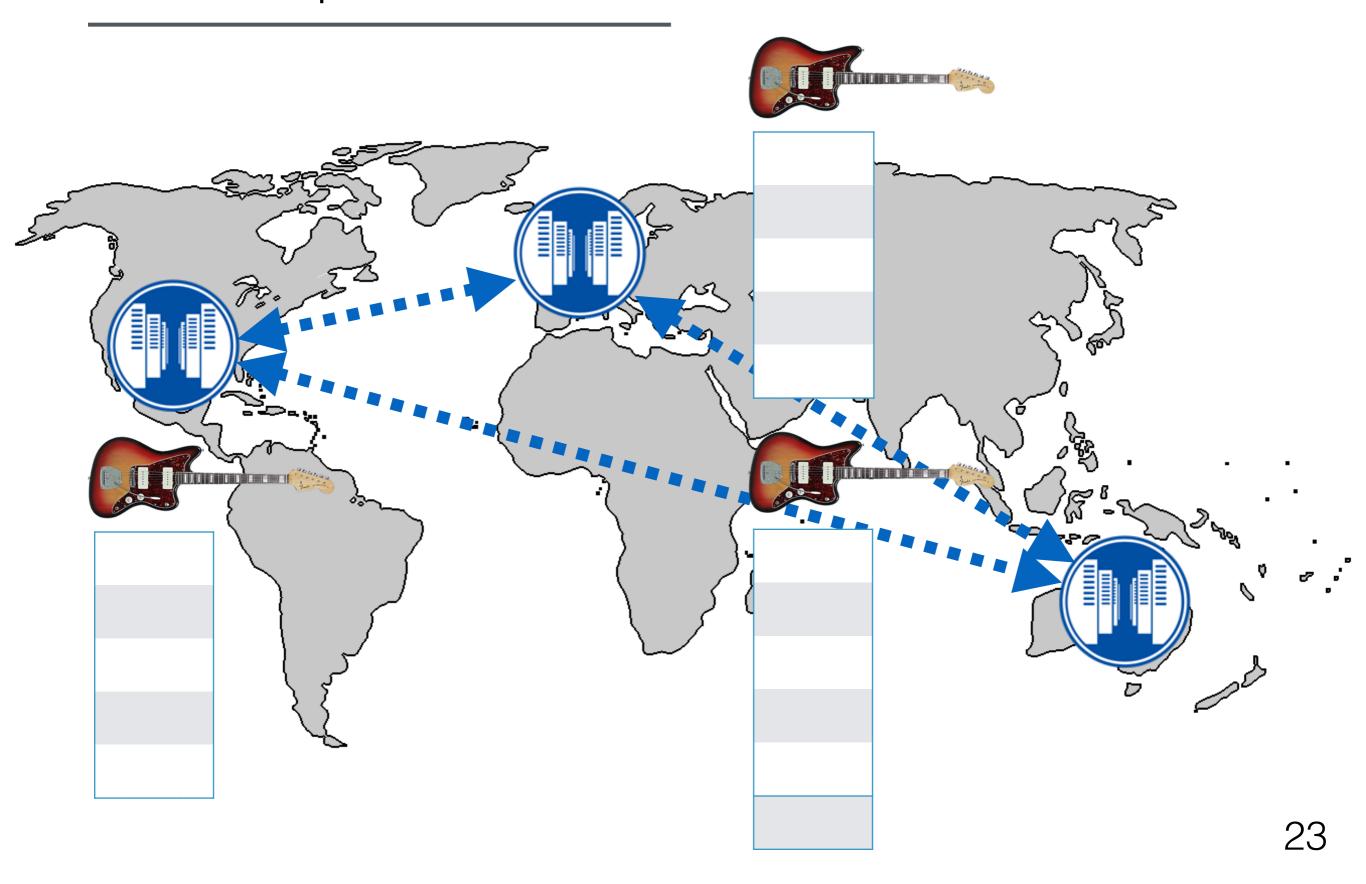
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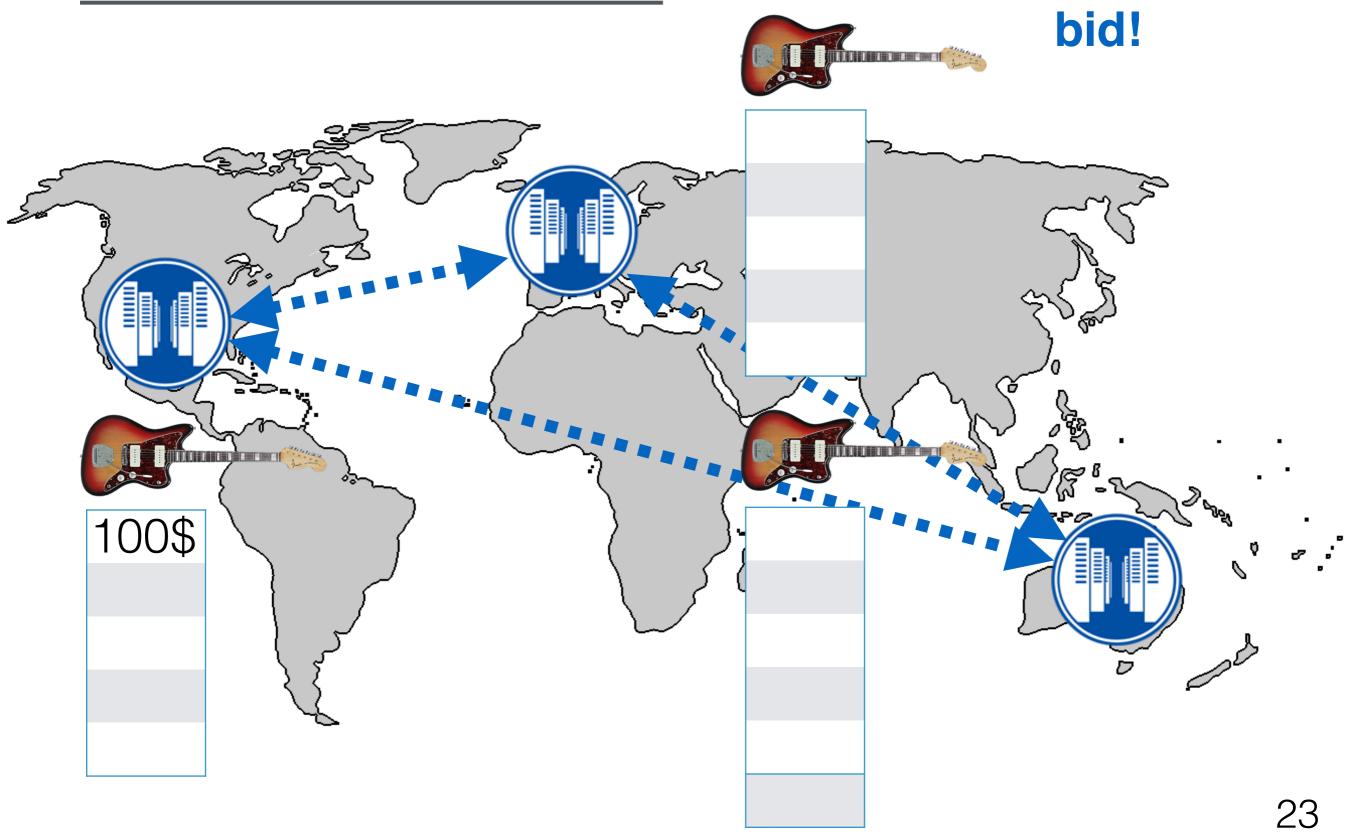
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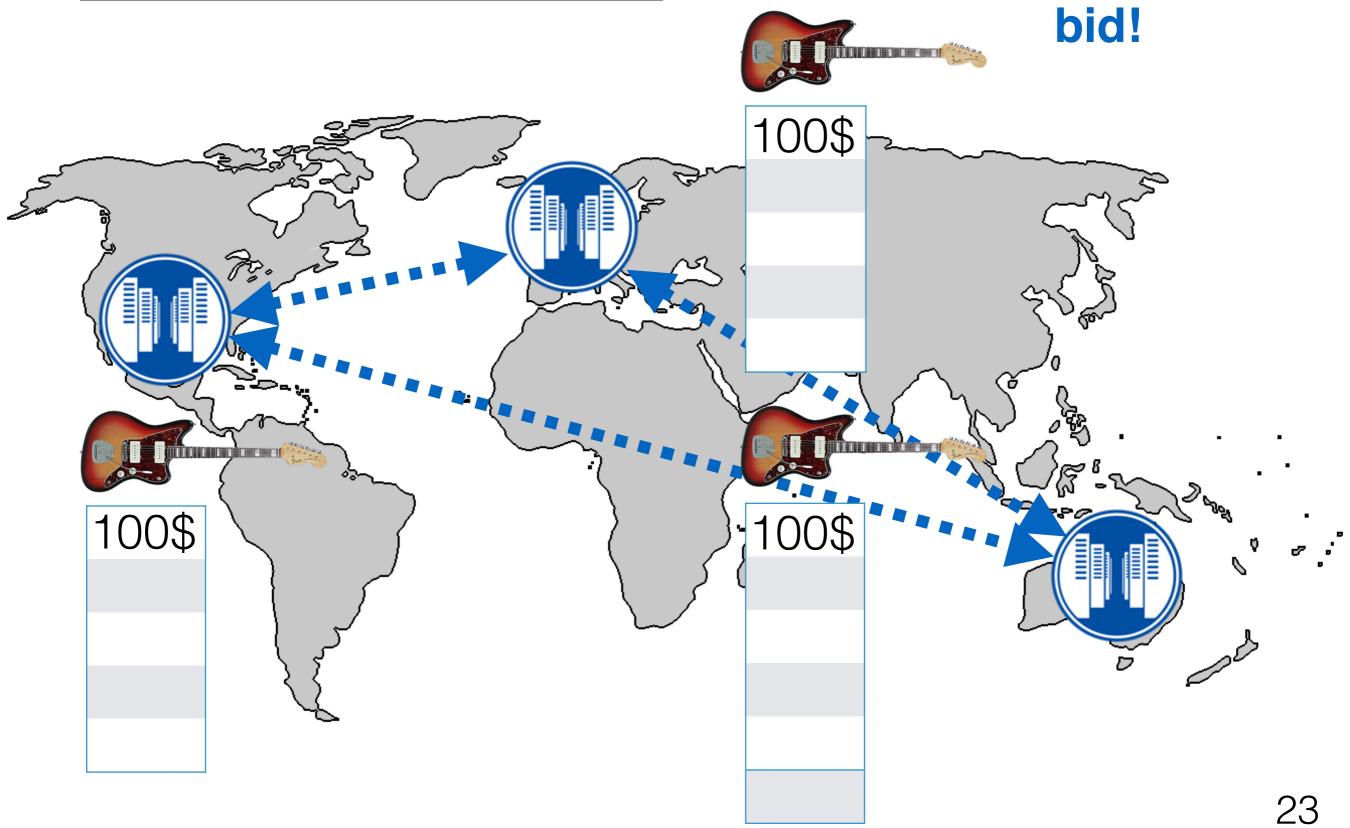




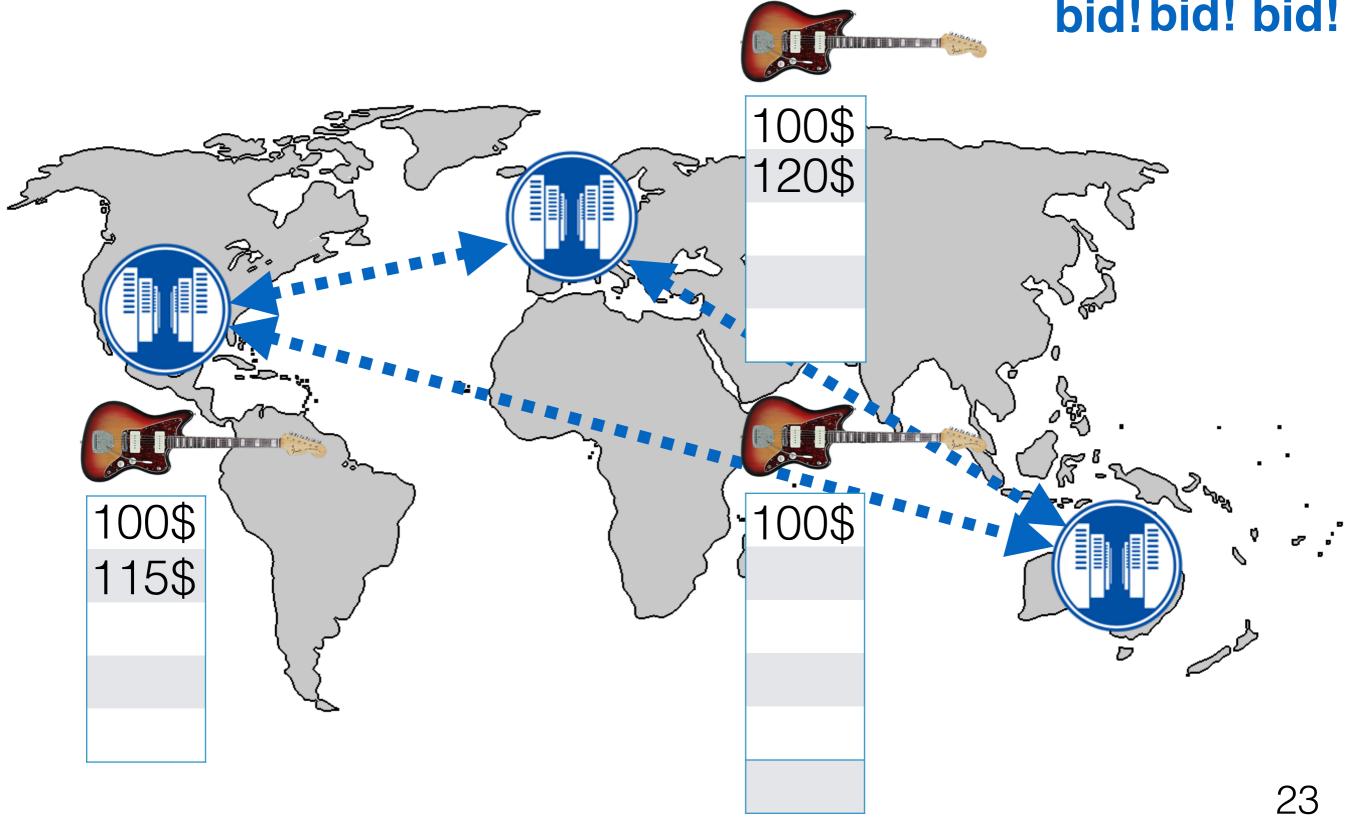




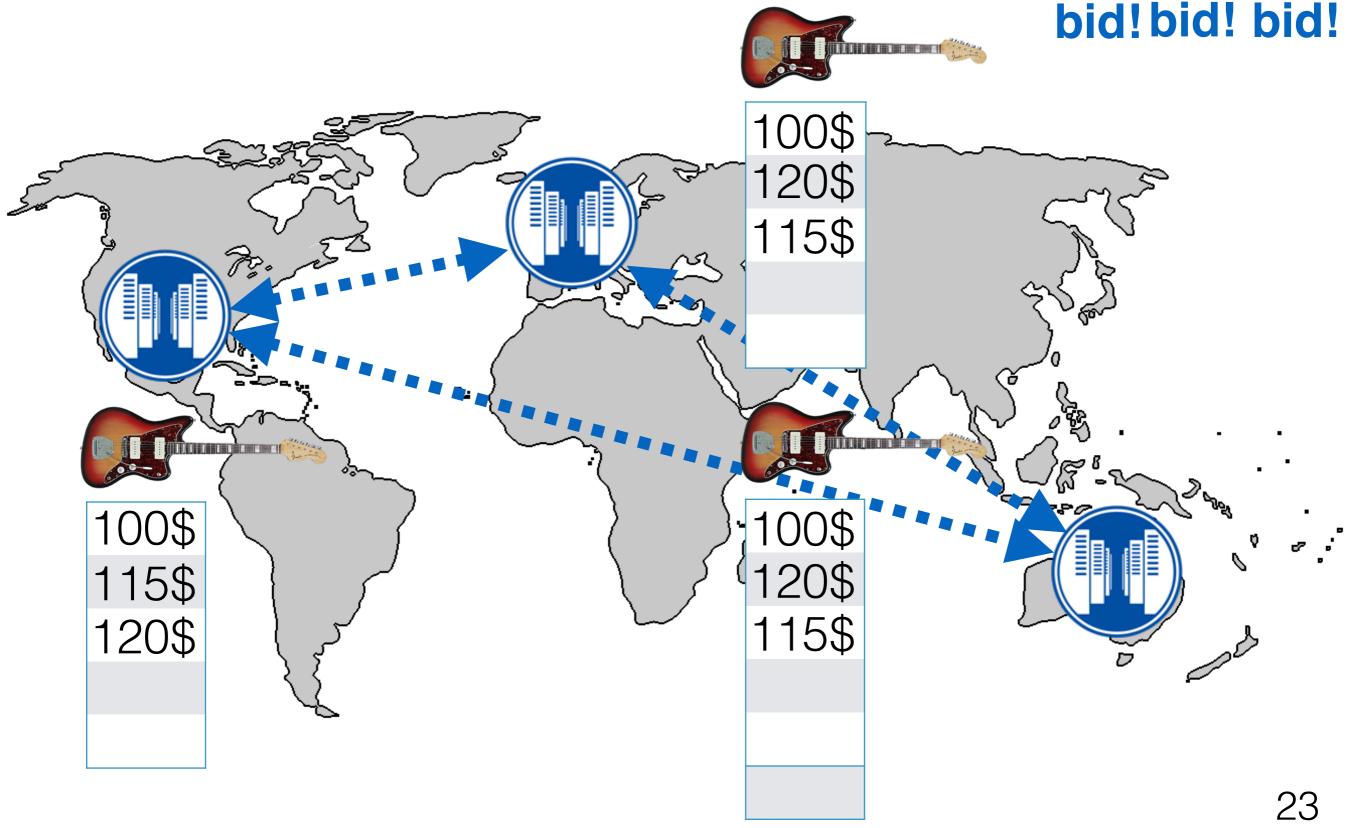




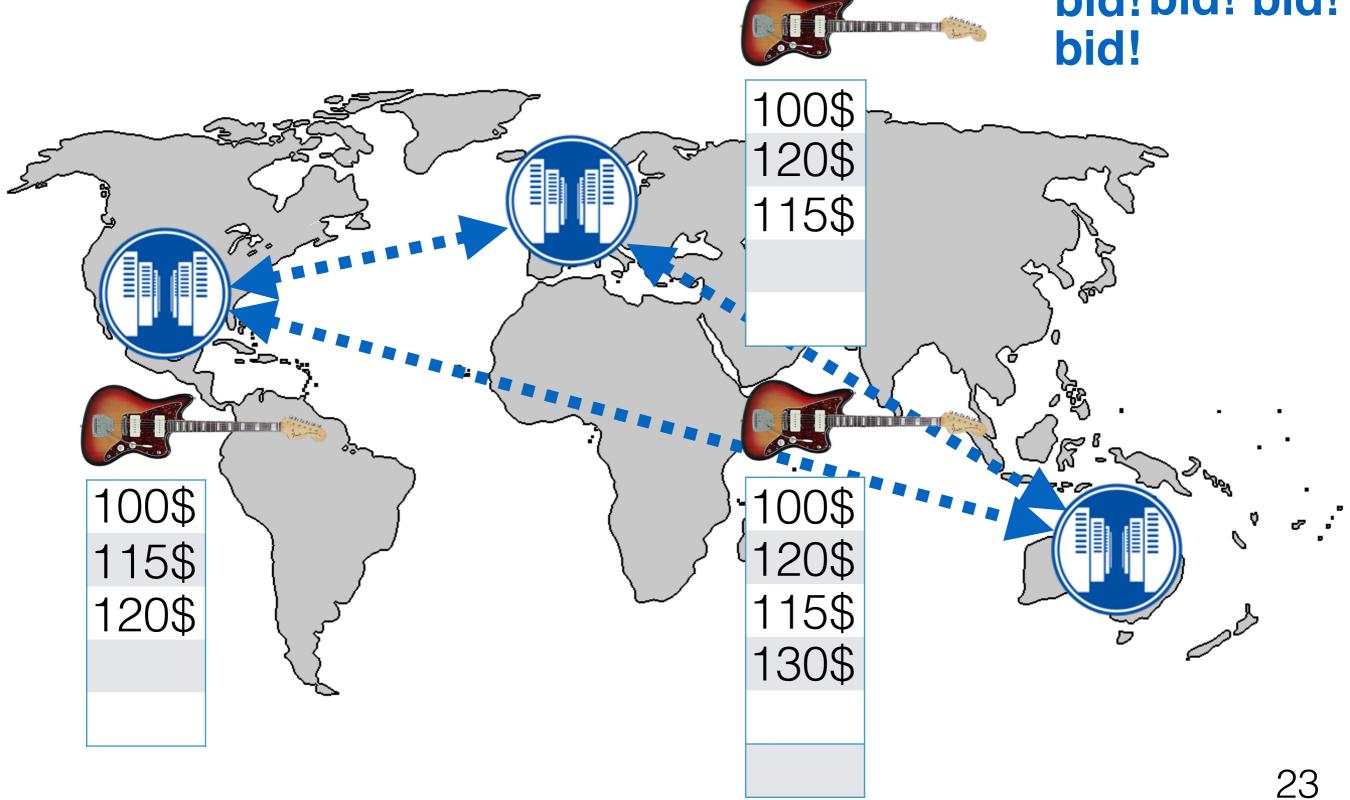
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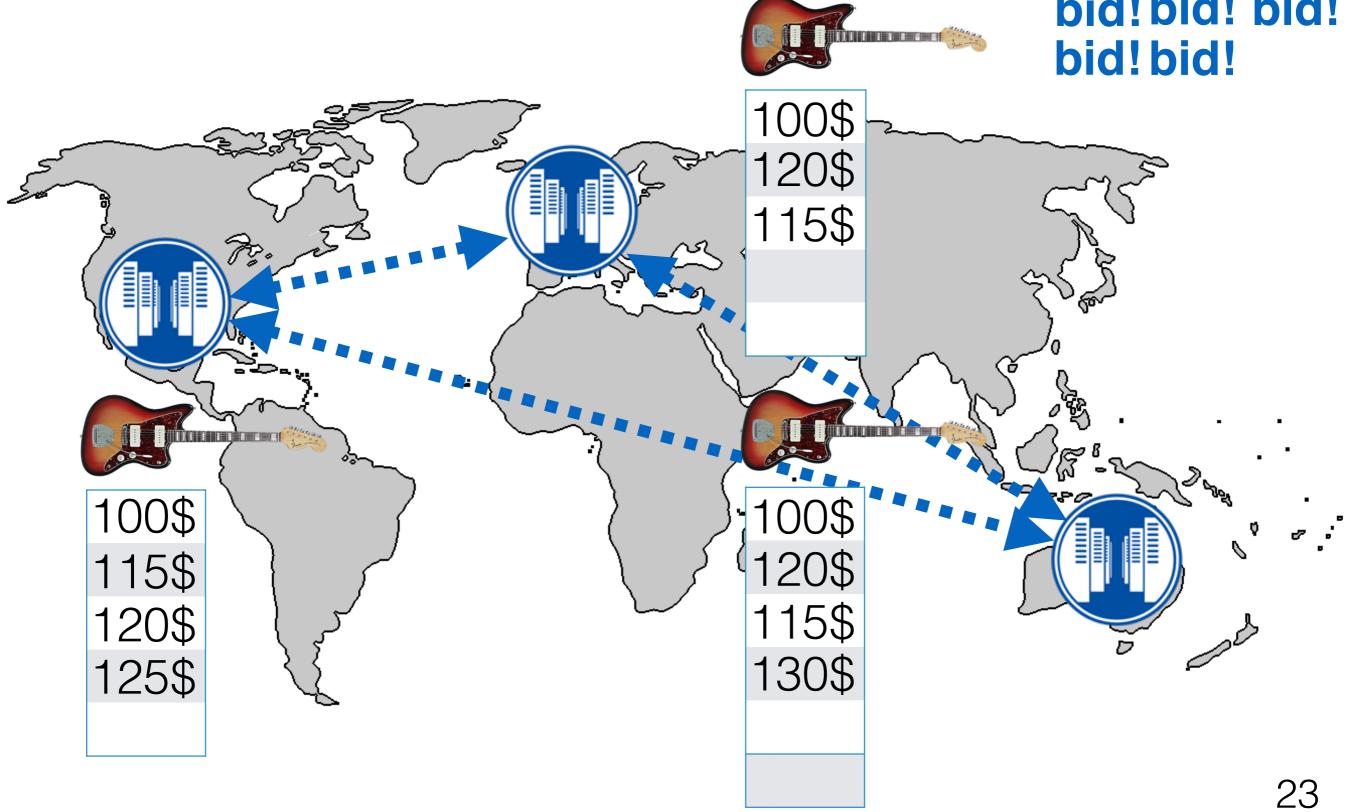
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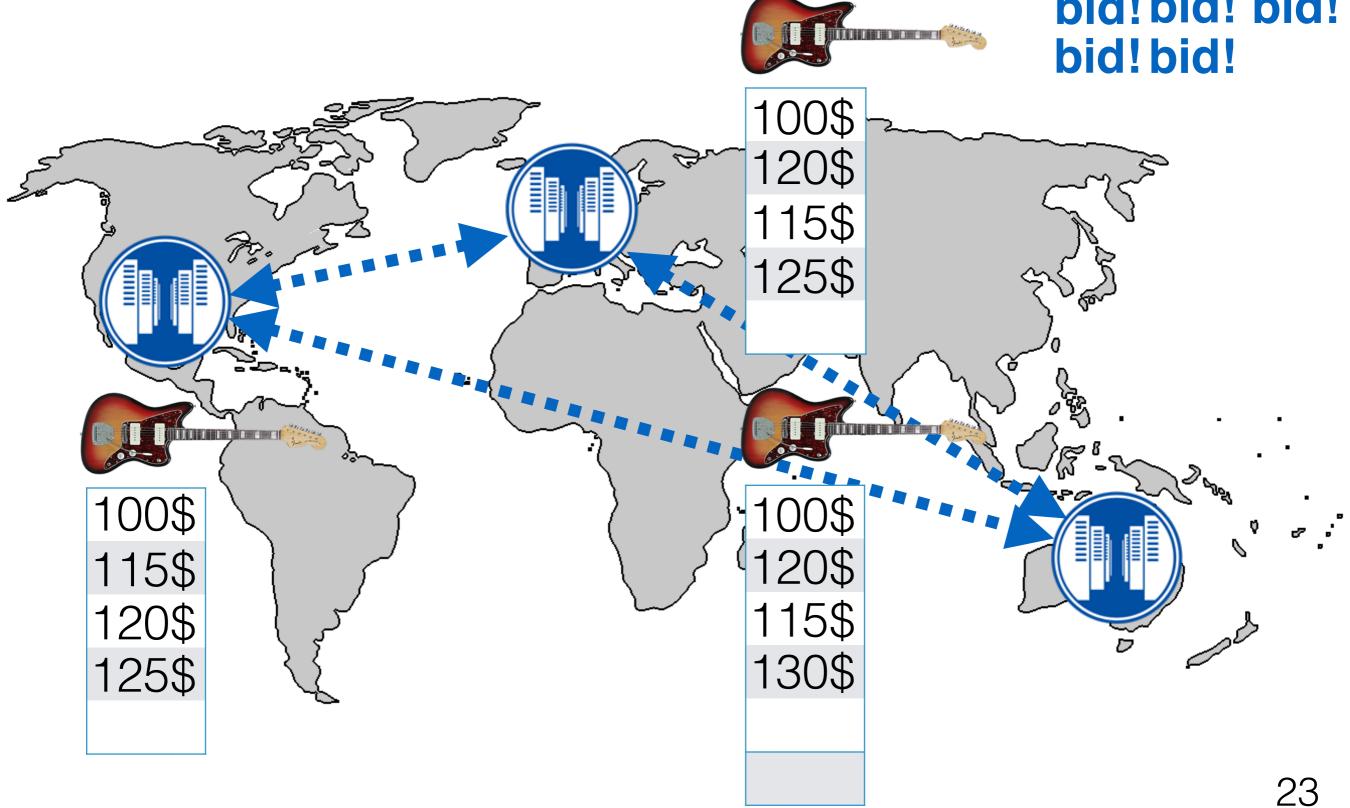
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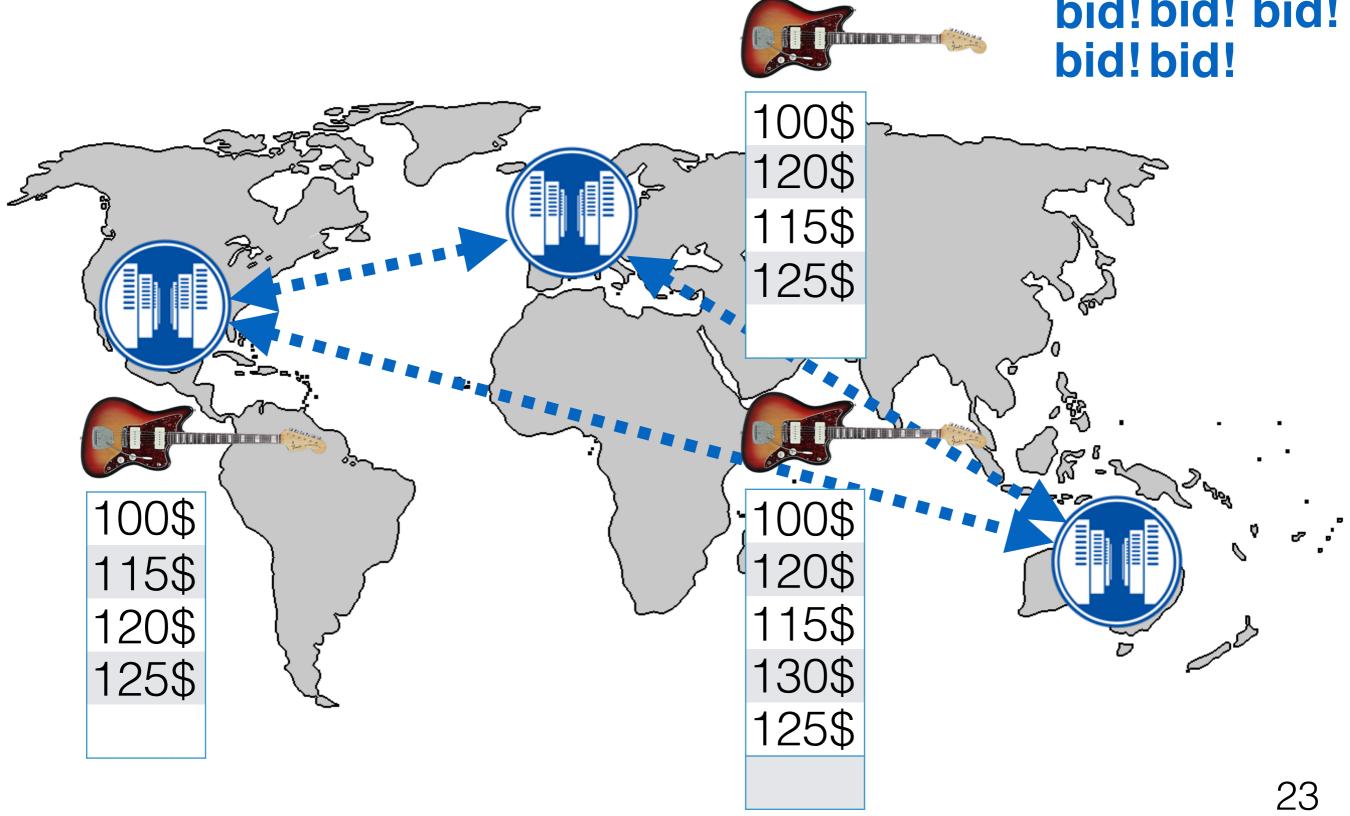
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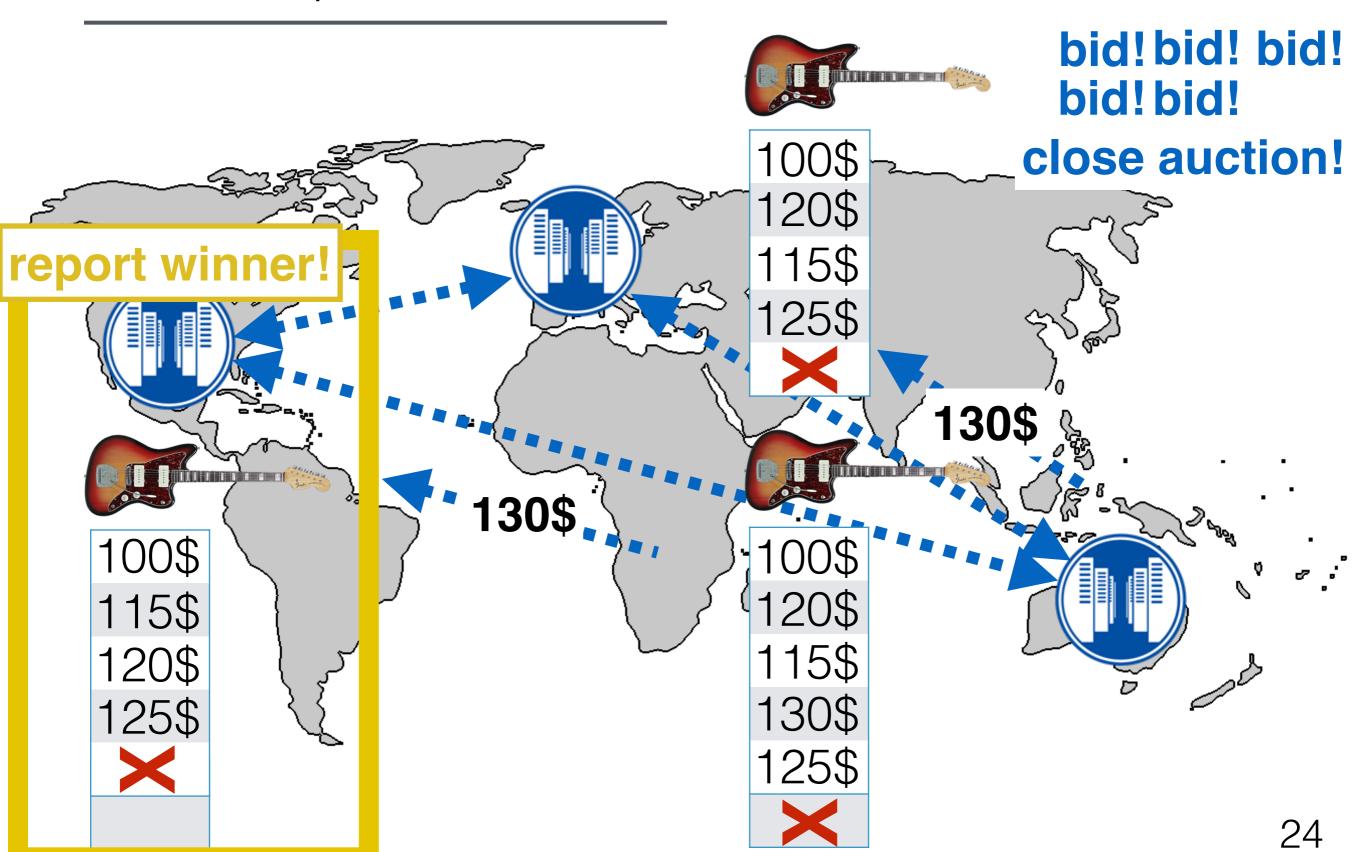
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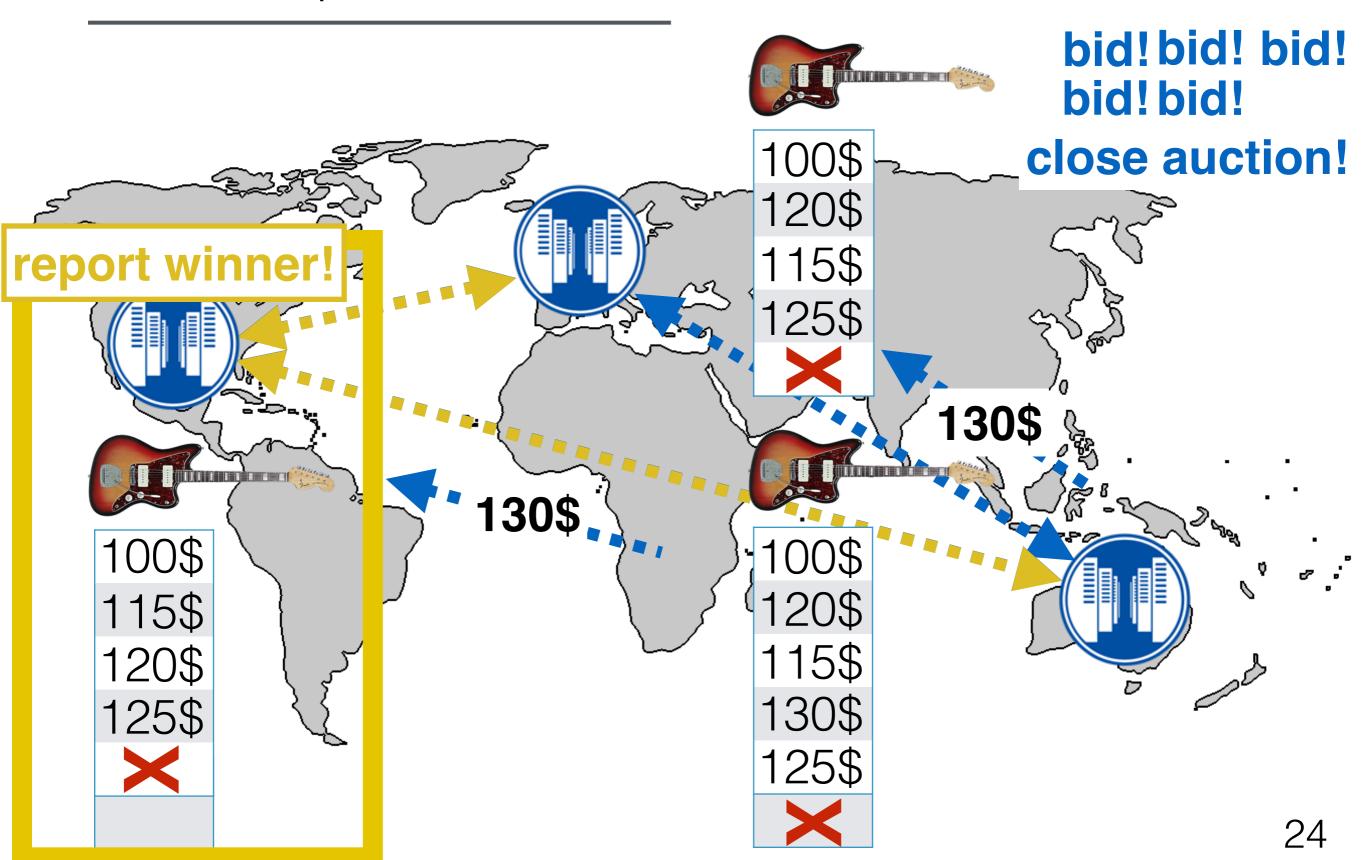
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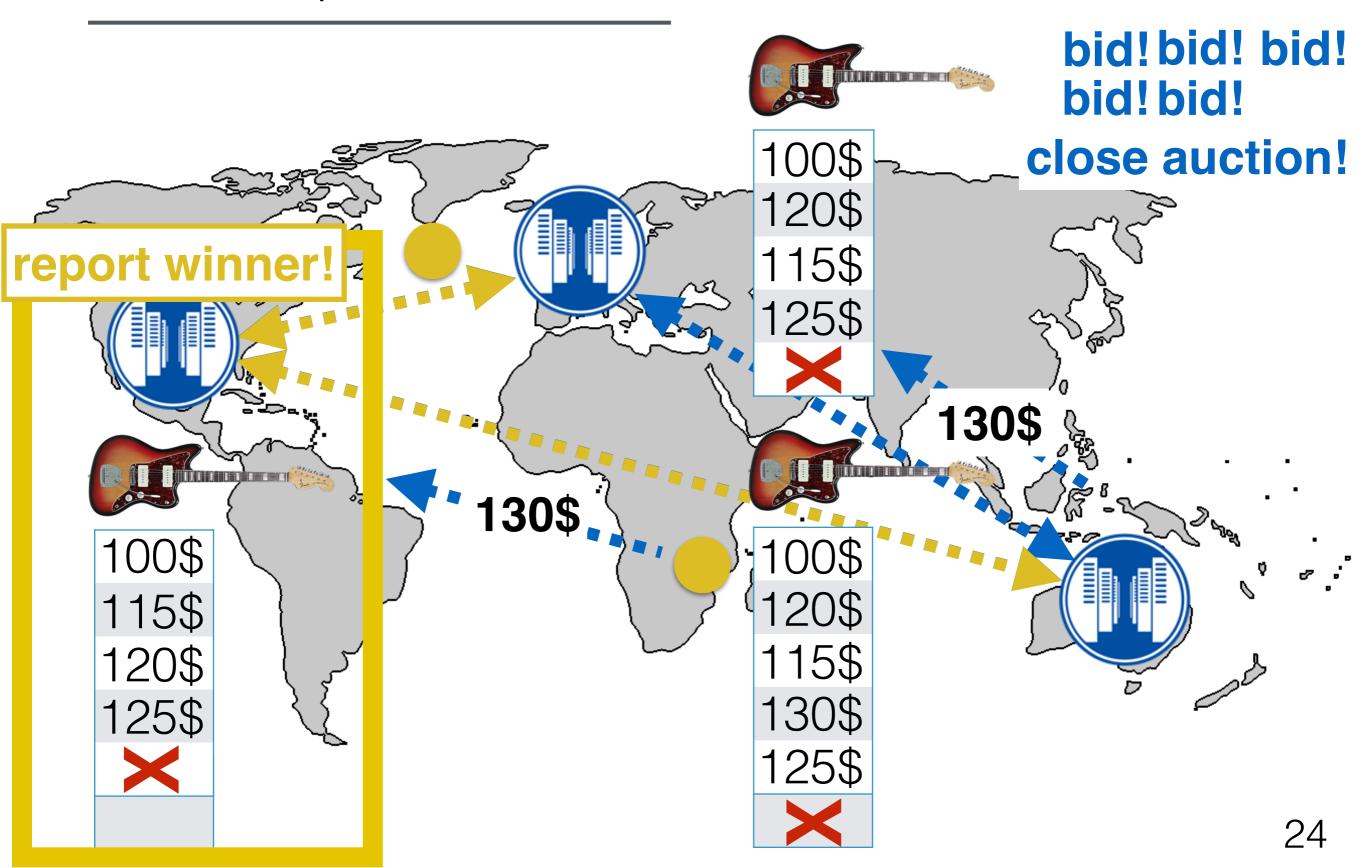
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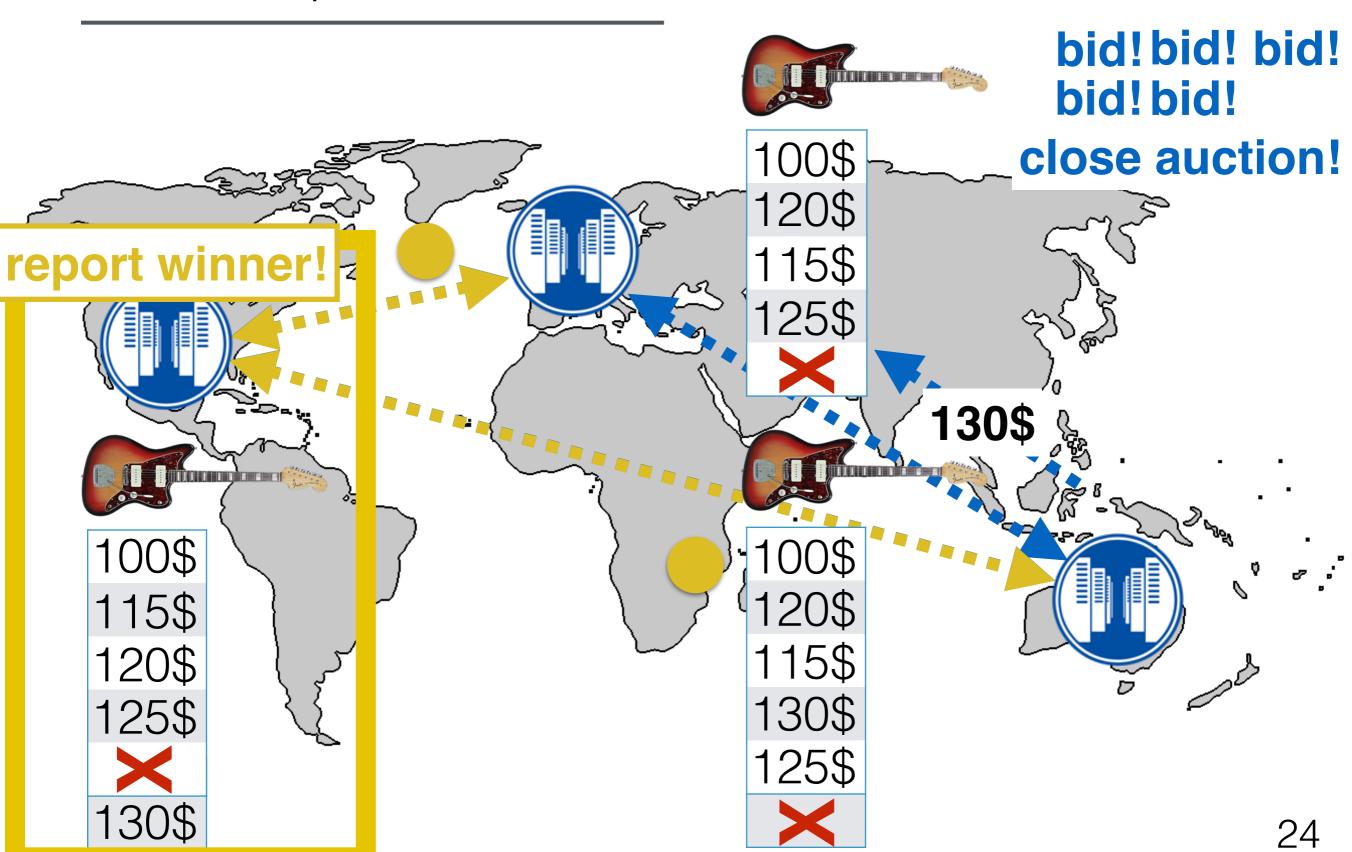
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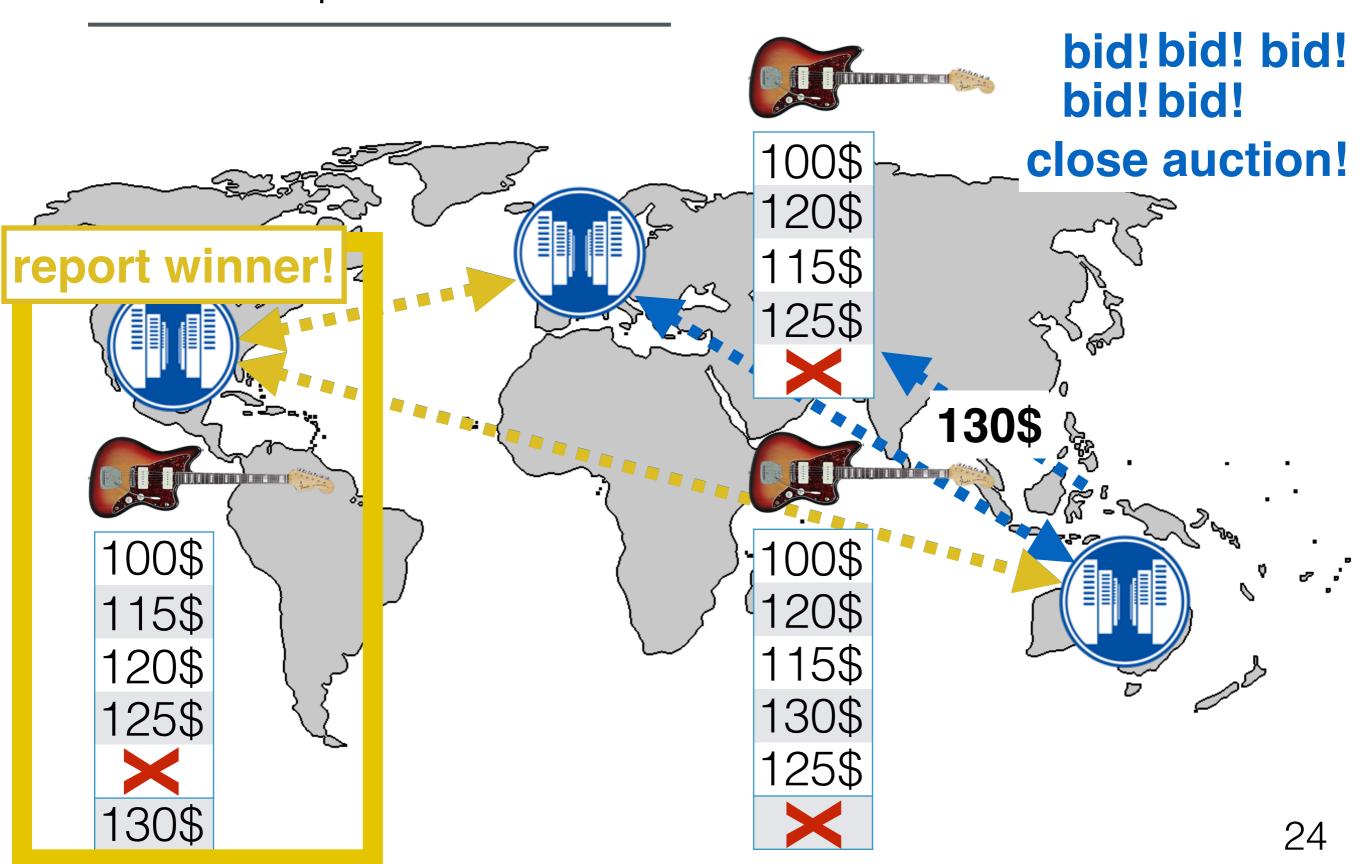
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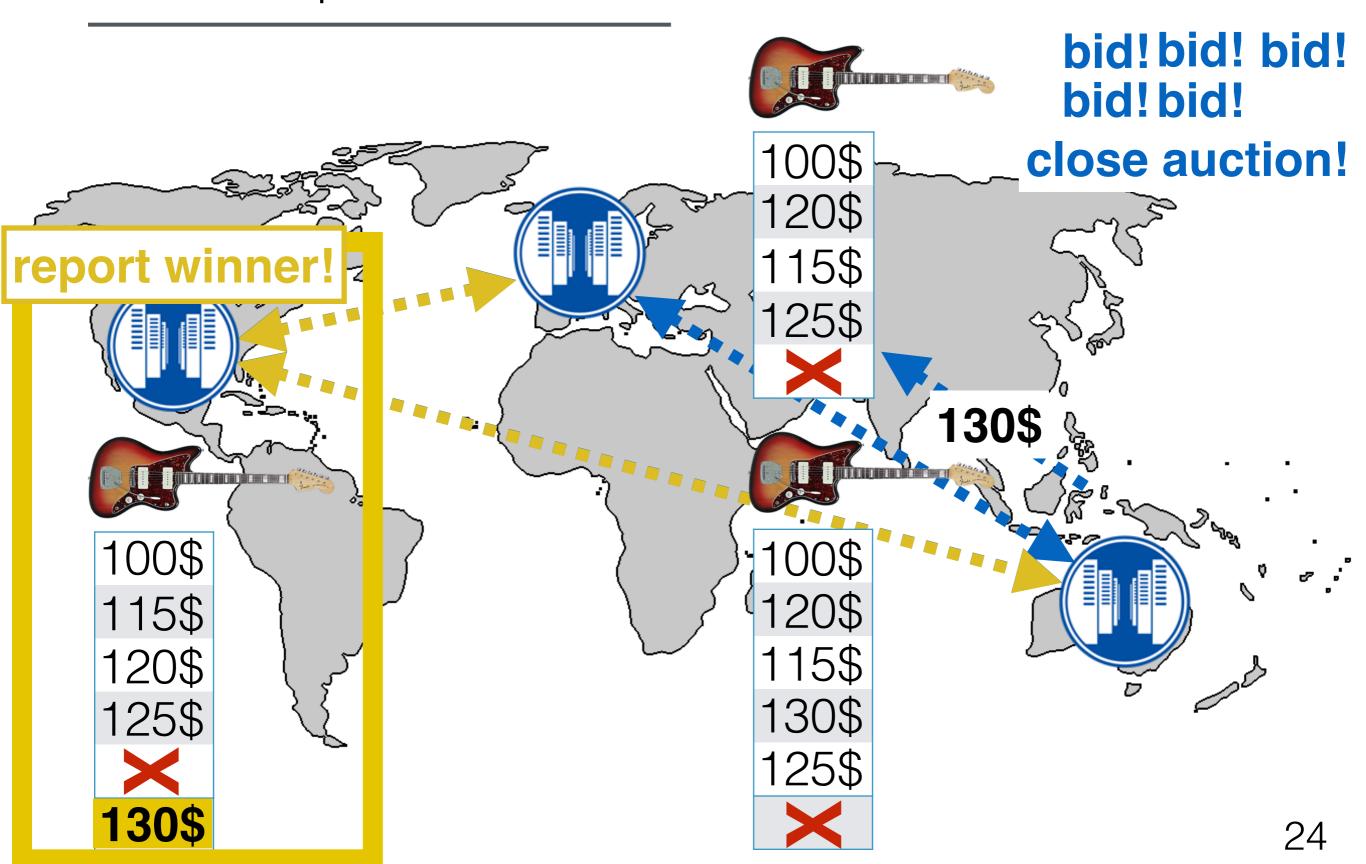


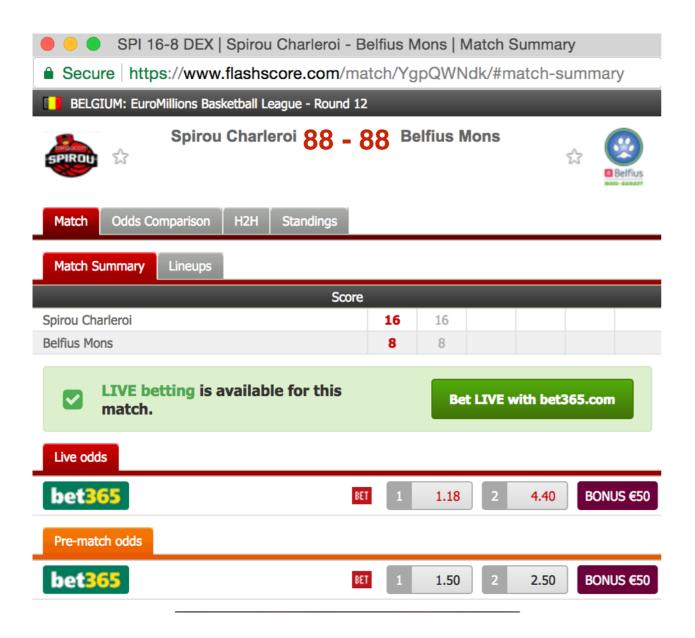


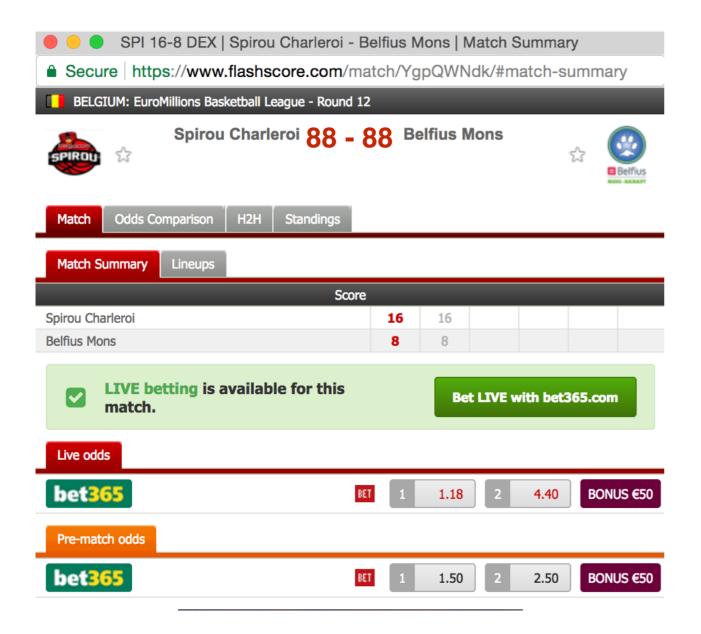




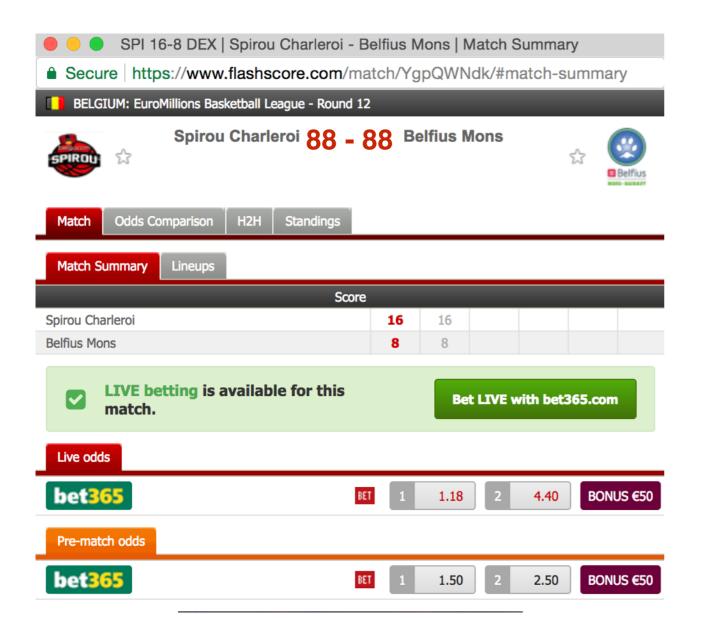






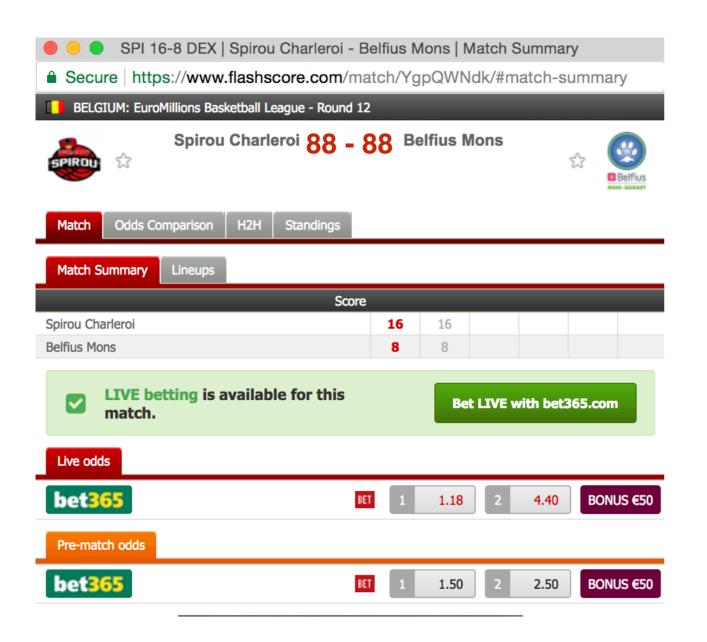


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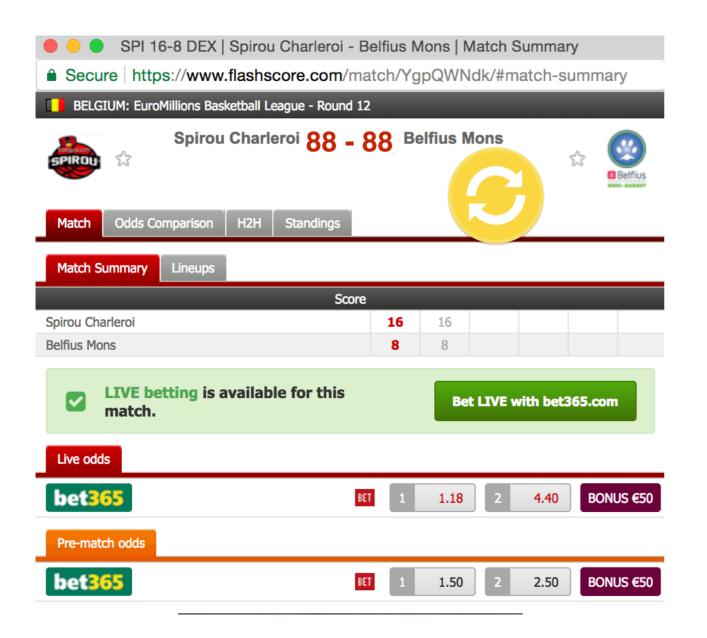
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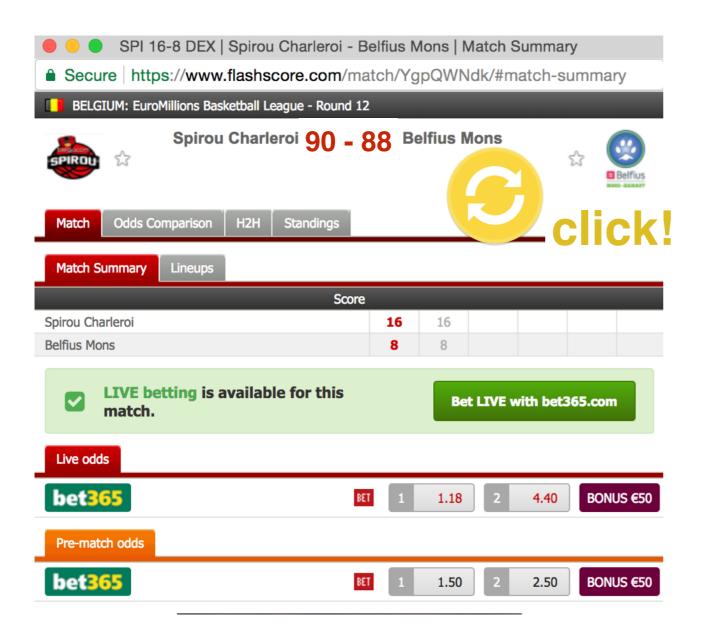
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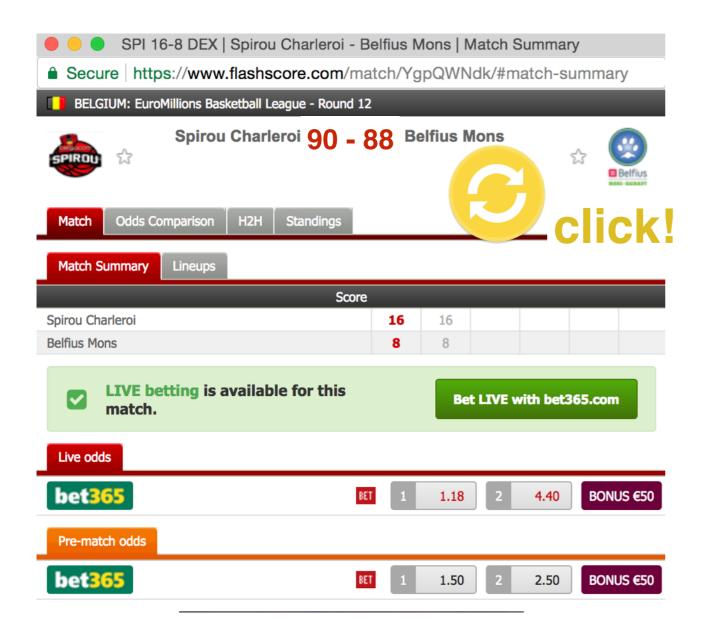
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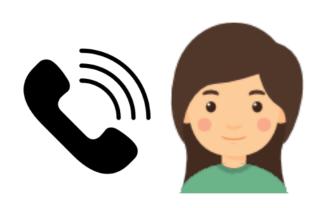
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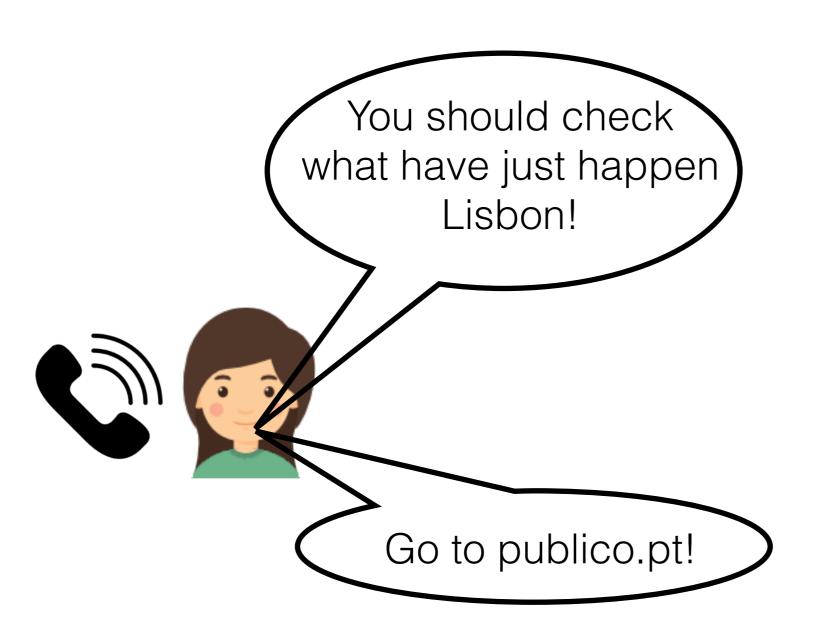
Useful to **read timesensitive** information



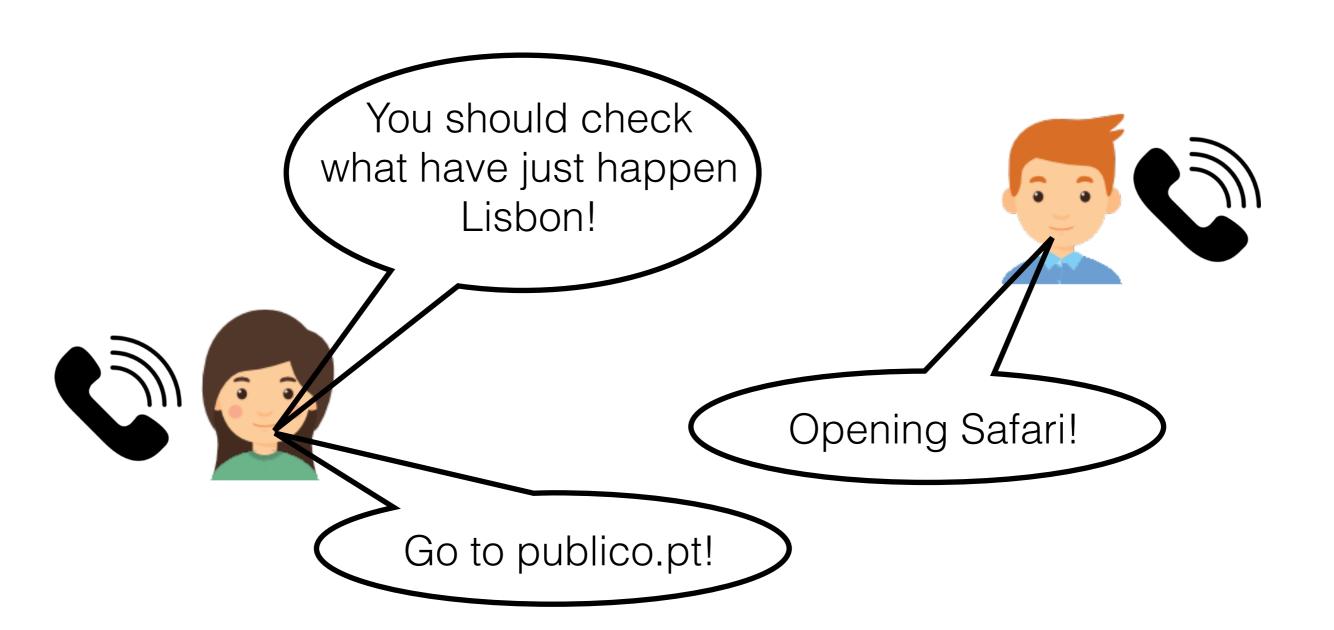


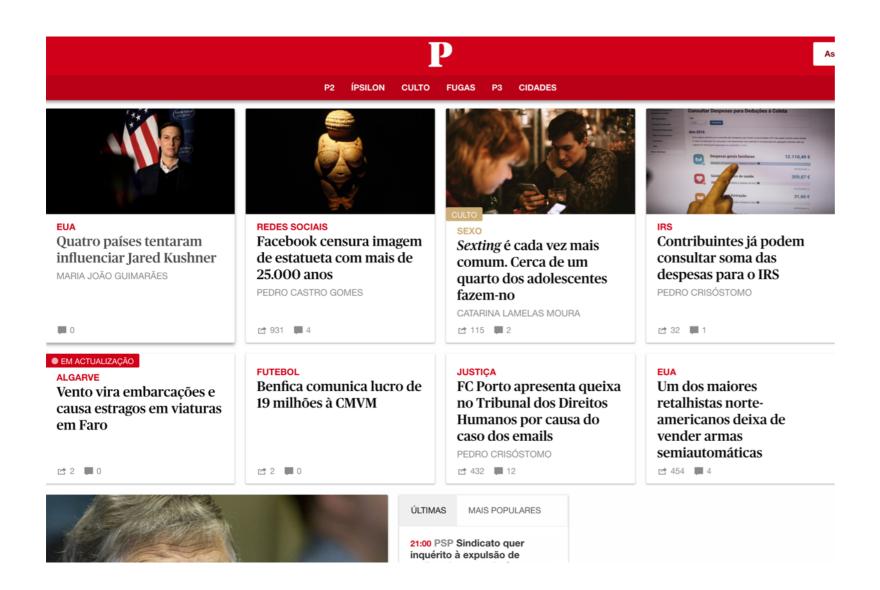




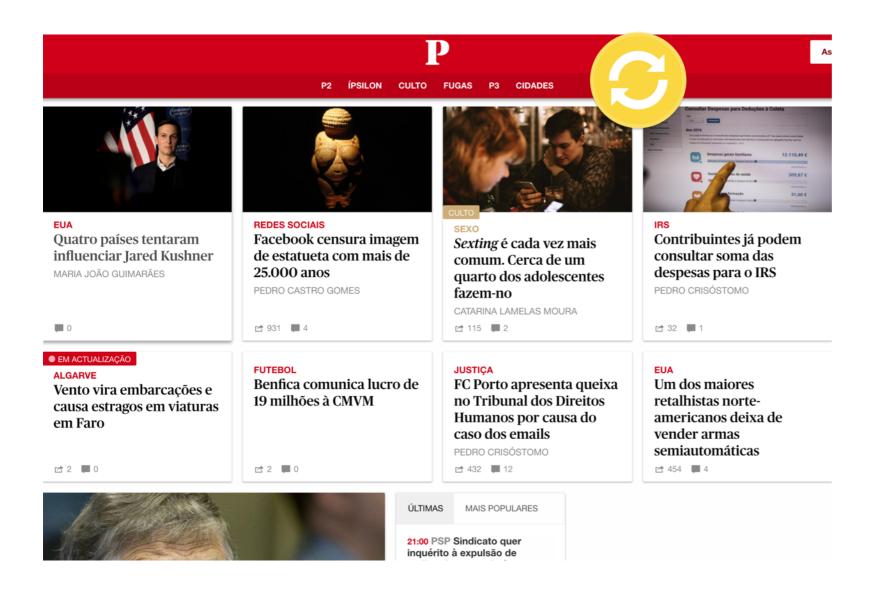




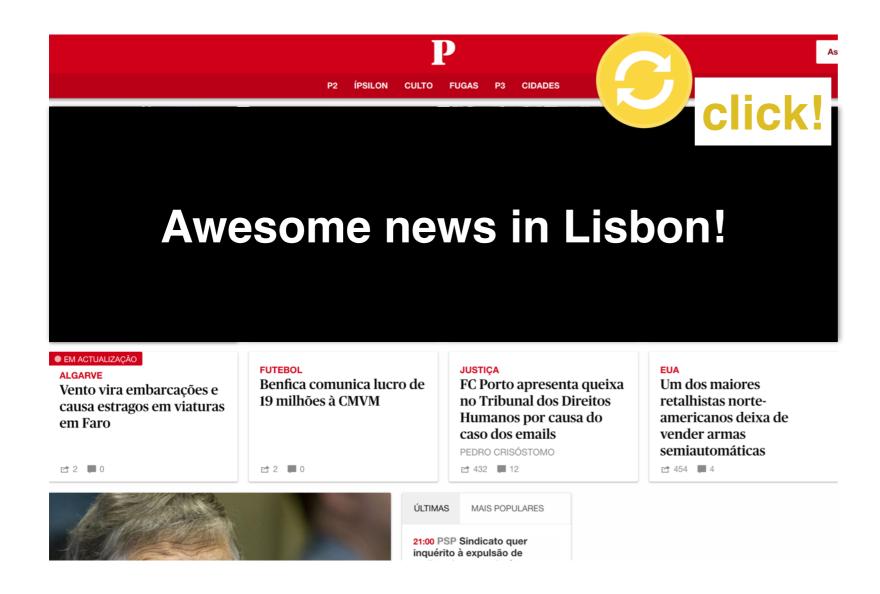




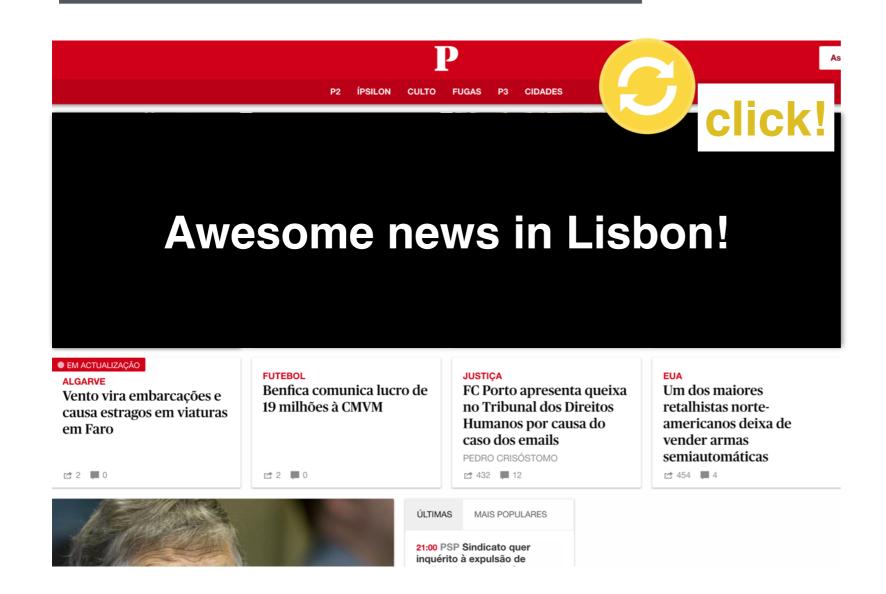














Useful to cope, at least partially, with back-channelling

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thus, the prepare phase of external operations does not require contacting all sites

bad things: we are making writes slightly more expensive

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good things: reads of internal operations still require a single site

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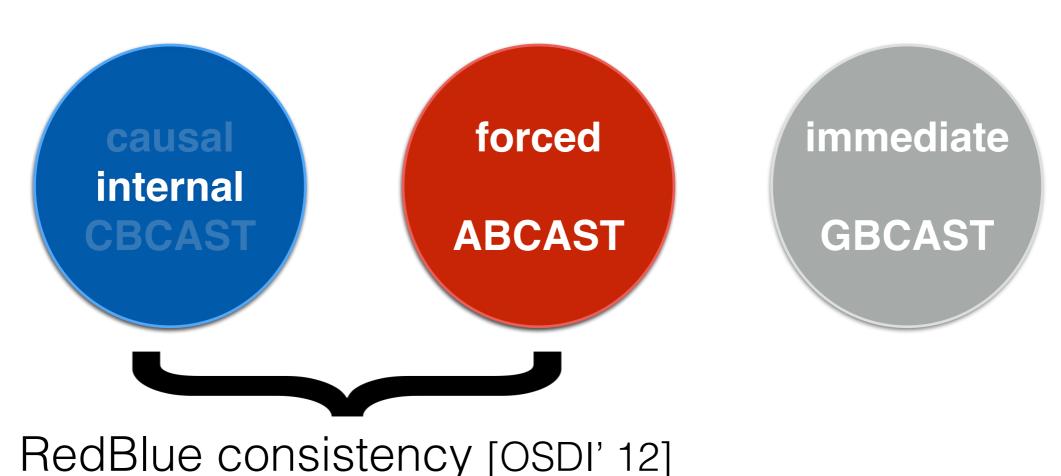
good things: reads of internal operations still require a single site

more good things: we enhance availability when compared to causal consistency, which otherwise is sticky available

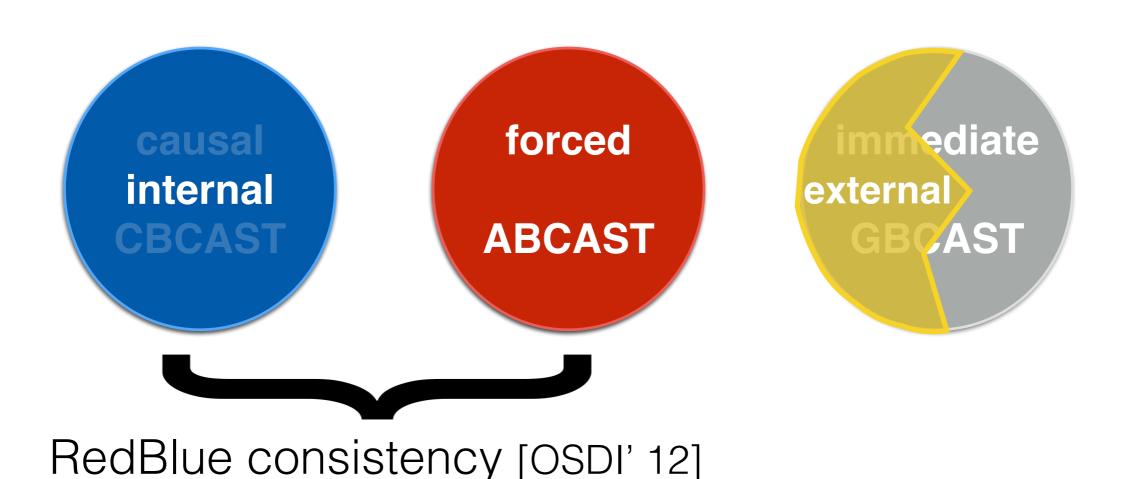








# closest to our work: lazy replication [TOCS' 92] and the ISIS work [TOCS' 87]



31





Performed at all replicas in the same order relative to all operations



Performed at all replicas in the same order relative to all operations

Ordered consistently with external events



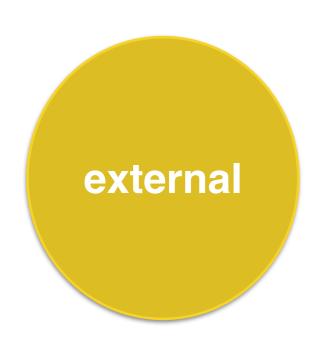
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Performed at all replicas in causal order

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Performed at all replicas in causal order

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shall we add a new operation that becomes all-present after completion but it is not order consistently with external events?

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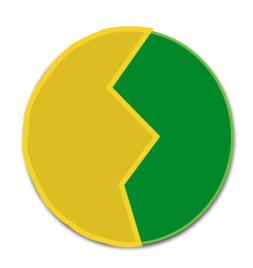
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Performed at all replicas in causal order

Ordered consistently with external events



The operation becomes allpresent after completion

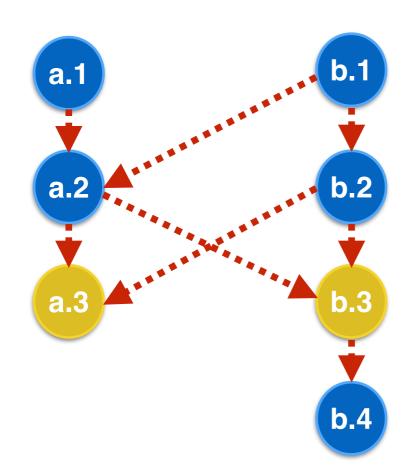
this could further help to beat backchanneling

shall we **add** mechanisms to ensure **convergence** (CRDTs) and to preserve invariants (forced/**red operations**) in order to build a complete system?

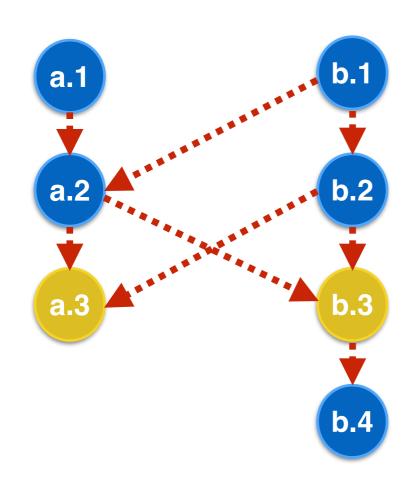
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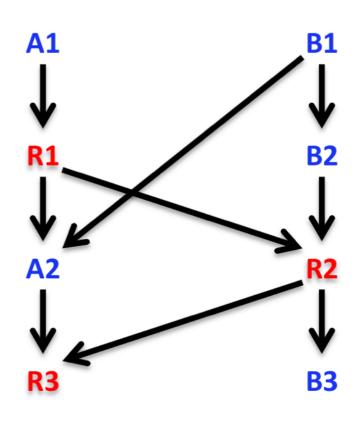
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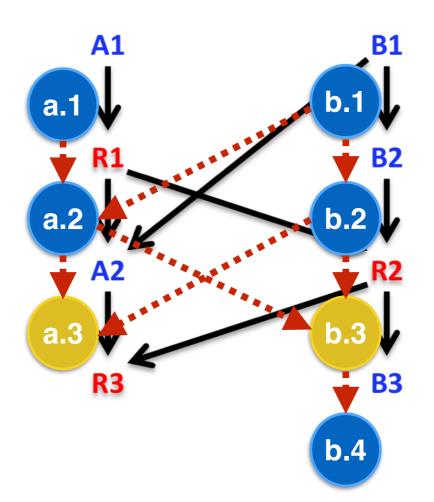


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how can we **efficiently implement** external causality in a geo-replicated, distributed system?

## Thanks! Questions?

we present a new consistency model: external causality

addresses the inherent tension between performance and meaningful semantics

combines causal consistency (strongest available) and external consistency (strongest semantically)